


Revisiting (R)CCA Security and Replay Protection

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Abstract. This paper takes a fresh approach to systematically characterizing, comparing, and understanding CCA-type security definitions for public-key encryption (PKE), a topic with a long history. The justification for a concrete security definition X is relative to a benchmark application (e.g. confidential communication): Does the use of a PKE scheme satisfying X imply the security of the application? Because unnecessarily strong definitions may lead to unnecessarily inefficient schemes or unnecessarily strong computational assumptions, security definitions should be as weak as possible, i.e. as close as possible to (but above) the benchmark. Understanding the hierarchy of security definitions, partially ordered by the implication (i.e. at least as strong) relation, is hence important, as is placing the relevant applications as benchmark levels within the hierarchy.

CCA-2 security is apparently the strongest notion, but because it is arguably too strong, Canetti, Krawczyk, and Nielsen (Crypto 2003) proposed the relaxed notions of *Replayable CCA* security (RCCA) as perhaps the weakest meaningful definition, and they investigated the space between CCA and RCCA security by proposing two versions of *Detectable RCCA* (d-RCCA) security which are meant to ensure that replays of ciphertexts are either publicly or secretly detectable (and hence preventable).

The contributions of this paper are three-fold. First, following the work of Coretti, Maurer, and Tackmann (Asiacrypt 2013), we formalize the three benchmark applications of PKE that serve as the natural motivation for security notions, namely the construction of certain types of (possibly replay-protected) confidential channels (from an insecure and an authenticated communication channel). Second, we prove that RCCA does not achieve the confidentiality benchmark and, contrary to previous belief, that the proposed d-RCCA notions are not even relaxations of CCA-2 security. Third, we propose the natural security notions corresponding to the three benchmarks: an appropriately strengthened version of RCCA to ensure confidentiality, as well as two notions for capturing public and secret replay detectability.

1 Introduction

1.1 Motivation

When designing a cryptographic security notion, it is of central importance to keep in mind the purpose and applications it is developed for. For CCA-2 secure encryption schemes³, the most important historical application is to enable confidential communication: assuming an insecure channel from Alice to Bob (over which ciphertexts are sent), and an authenticated channel from Bob to Alice (over which the public key can be transmitted authentically), the scheme should construct a confidential channel, i.e. an idealized object with the property that whatever Alice sends to Bob does not leak any information to an attacker (except possibly the length of the message), and where the only active capability of the attacker is to inject new messages (uncorrelated to Alice’s inputs)⁴. Coretti, Maurer, and Tackmann [9] proved that indeed CCA-2 security is sufficient for this construction to be achieved, by having Bob generating a key-pair, sending the public key authentically to Alice, and by letting Alice encrypt all messages with respect to the obtained public key. It is also known that CCA-2 security is actually too strong for this task: a CCA-2 secure scheme can be easily modified, for example by appending a single bit to ciphertexts which is ignored by the decryption algorithm, to yield a scheme that is not CCA-2 secure but still allows to achieve a confidential channel.

To address the question what weaker security notion(s) would actually match more closely to the application of secure communication, Canetti, Krawczyk, and Nielsen [7] initiated the study of relaxing CCA-2 security. They formalize an entire spectrum of security notions: at the weakest end, they propose RCCA security, which for large message spaces (size super-polynomial in the security parameter) is known to achieve confidential channels [9]. This fact has bolstered RCCA security into becoming the default security notion in settings where CCA-2 is not achievable, such as in rerandomizable encryption schemes [22,14] and updatable encryption schemes [16]. Intuitively, a scheme can be RCCA secure even if it is easy to create from a known ciphertext another one that still decrypts to the same message. Inheriting from prior work on relaxing CCA-2 security, most notably [1,24,17], they further provide formalizations for intermediate notions between CCA-2 and RCCA. These so-called detectable notions of RCCA security further demand that modifications of an already known ciphertext can be efficiently detected—either with the help of the secret key (*sd-RCCA*) or the public key only (*pd-RCCA*) yielding two separate security notions. These notions of detectable RCCA security, and in particular *pd-RCCA*, are designed to capture an appealing property of CCA-2 security, namely that replays can be efficiently filtered out, which is desirable in applications like voting or access-control encryption, where a trusted third party must perform the filtering without access to the secret key [3].

³ Note that throughout this work, if not otherwise stated, we refer to the indistinguishability-based versions of security notions.

⁴ Hence, the confidential channel does not provide any authenticity to Bob.

It has however never been formally investigated whether the detectable notions are suitable to capture the security of the intended application of replay detection. Moreover, our analysis shows that these detectable RCCA notions (i.e. *pd-RCCA* and *sd-RCCA*) are actually not proper relaxations of *CCA-2*, in that they are not implied by *CCA-2*.

In this work, we fill this gap and provide a systematic treatment of these relaxations of *CCA-2* security using the Constructive Cryptography framework by Maurer and Renner [19,18] and building upon the work of Coretti et al. [9]. We formalize the intuitive security goals that *RCCA* security and the detectable *RCCA* security notions aim to achieve, yielding what we call benchmarks to assess whether the existing security notions are adequate. We observe that none of the previous notions seems to allow a proof that they meet this level of security and therefore propose new security notions for detectable *RCCA* security (which can be regarded as the corrections of the existing ones), show which benchmarks they achieve, and prove that they are implied by *CCA-2*. In summary, this shows that the newly introduced notions are placed correctly in the spectrum between *CCA-2* and *RCCA* and that they can be safely used in the intended applications.

1.2 Overview of Contributions

A systematic approach to *RCCA* and replay protection. Following the constructive paradigm, a construction consists of three elements: the assumed resources (such as an insecure communication channel), the constructed or ideal resource (such as a confidential channel), and the real-world protocol. A protocol is said to achieve the construction, if there is a simulator such that the real world (consisting of the protocol running with the assumed resources) is indistinguishable from the ideal system (consisting of the ideal resource and the simulator). This way, it is ensured that any attack on the real system can be translated into an attack to the ideal system, the latter being secure by definition.

Building upon the work of Coretti et al. [9], we present three benchmarks to approach the intended security of *RCCA* and replay protection:

- The construction of a confidential channel between Alice and Bob from an insecure communication channel (and an authenticated channel to distribute the public key). This is arguably the most natural goal of confidential (and non-malleable) communication. An encryption scheme should achieve this construction by having Bob generating the key-pair and sending the public key to Alice over the authenticated channel. Alice sends encryptions of the messages over the insecure channel to Bob, who can decrypt the ciphertexts and output the resulting messages. This benchmark is formalized in Sect. 3.1.
- The construction of a replay-protected confidential channel from (essentially) the same resources as above. A replay-protected confidential channel is a channel that only allows an attacker to deliver each message sent by Alice at most once to Bob. This construction captures the most basic form of replay protection. An encryption scheme can be applied as above, except that Bob must make use of the secret key (and a memory resource to store received

ciphertexts) to detect and filter out replays. This construction is formalized in Sect. 3.2.

- The construction of a replay-protected confidential channel from basically the same resources, but where the task of detecting replays is done by a third-party, say Charlie, that does not need to have access to Bob’s secret key. Hence, an encryption scheme is employed as above, but the task of filtering and detecting replays can be outsourced to any party possessing the public key (having sufficient memory to store the received ciphertexts). This benchmark is formalized in Sect. 3.3.

We note that only the first benchmark is taken from existing literature [9] (which is an abstract version of the UC-formalization $\mathcal{F}_{\text{MPKE}}$ defined in [7])⁵, while the other benchmarks are new formulations and variants of the known goal of replay protection. The benefits of all our benchmarks is that they yield a precise way to assess the guarantees provided by a security notion for an encryption scheme: does a scheme secure with respect to a certain notion achieve the above construction(s)?

New intermediate notions between CCA-2 and RCCA. We propose three game-based security notions, each designed to suffice for achieving the intended benchmark. The abbreviations stand for confidential (cl), secret-key replay protection (srp), and public-key replay protection (prp):

- We first propose IND-cl-RCCA, a security notion which is sufficient to achieve confidential communication even for small message spaces, which we prove in Sect. 5.1. This is the weakest new notion we introduce and we prove that it achieves the first benchmark; cl-RCCA should then take the role of RCCA as the default security notion when one aims at the design of schemes that enable confidential communication (in particular when the message space size is small). Note that cl-RCCA is strictly stronger than RCCA since the latter does not achieve confidential communication for small message spaces. For completeness, we give a proof of this fact in Appendix A.⁶
- The second security notion we introduce is IND-srp-RCCA and it achieves the second benchmark: realizing a replay protected confidential channel. The notion is hence designed to enable the implementation of a replay-protection mechanism by the receiver, who knows the secret decryption key. We also argue why the strengthening compared to cl-RCCA (and sd-RCCA) is needed to achieve replay-protection: from a conceptual perspective, implementing a replay-protector as part of the receiver requires the detection of replays without necessarily ever seeing the original ciphertext by the sender which

⁵ We note that all our results do not depend on the specific details of the underlying composable framework.

⁶ We note that NM-RCCA [7], which is stronger than IND-RCCA, does not seem to be sufficient to achieve the first benchmark either.

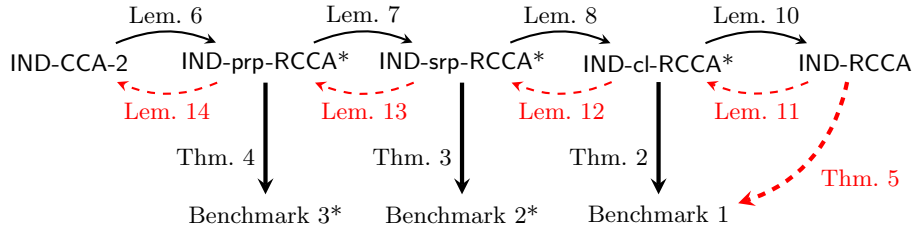


Fig. 1. New notions of security between CCA-2 and RCCA, and their relations to each other and to the benchmarks. Solid black arrows denote implications and dashed red arrows denote separations. The new security notions introduced in this paper are marked with *.

is a security requirement that is not captured by cl-RCCA (nor sd-RCCA).⁷ The notion and the construction proof appear in Sect. 5.2.

- We finally propose a security notion to capture the idea of publicly-detectable RCCA that we call IND-prp-RCCA. This notion is sufficient to achieve the third benchmark and therefore captures the outsourced replay-protection mechanism that was originally envisioned from pd-RCCA. This notion and the construction proof appear in Sect. 5.3.

We finally show that all these notions can be strictly separated: IND-RCCA security, the weakest notion considered in this work, is strictly weaker than IND-cl-RCCA. The latter is strictly weaker than IND-srp-RCCA, which is in turn strictly weaker than IND-prp-RCCA. Finally, IND-prp-RCCA is strictly weaker than IND-CCA-2 security. These results are proven in Sect. 6; Fig. 1 illustrates all these new notions, their relations to each other and to the benchmarks.

Technical inconsistencies with existing pd-RCCA and sd-RCCA notions.

Numerous weaker versions of CCA-2 security have been proposed [1,7,17,24] which are essentially equivalent versions of what is formalized in [7] as *publicly detectable* (pd)-RCCA and *secretly detectable* (sd)-RCCA. We show for the given formalizations that the notions are generally not implied by CCA-2 security (unless one would restrict, for example, explicitly to the case of deterministic decryption [1] and perfect correctness), which seems to be a rather unintended artifact of the concrete definition as we show in Sect. 4. While these shortcomings can be fixed, the existing notions do not appear to suffice to achieve the intended benchmarks for replay protection (see Sect. 5), leaving the state of affairs unclear, as depicted in Fig. 2. This justifies the need to propose new intermediate notions that provably avoid these shortcomings: on one hand, our notions are implied by CCA-2, and on the other hand, they deliver the desired level of security required by a replay protection mechanism. The security notions and results of

⁷ More concretely, the simulator in the construction proof of a confidential channel only requires the (much milder) detection of *honestly* generated ciphertext replays.

this paper clean up the space between CCA-2 and RCCA security, yielding, as aforementioned, a clean hierarchy of security notions as depicted in Fig. 1: not only all notions are separated, but also we show that each of the notions we introduce is sufficient for achieving each of the benchmarks.

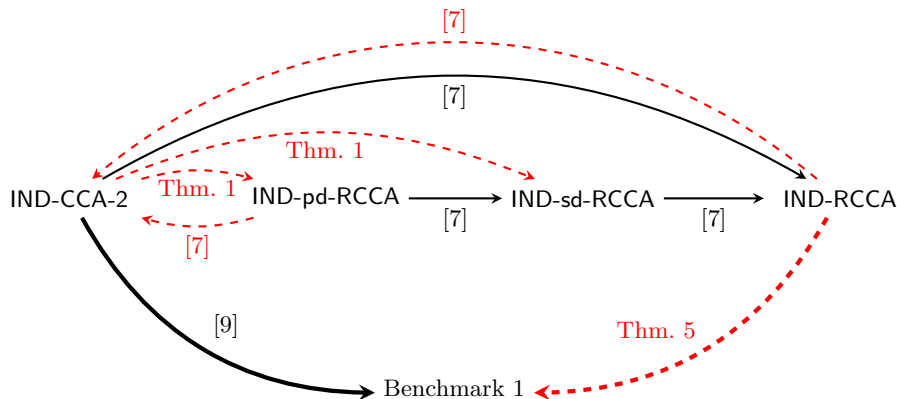


Fig. 2. Relations between the notions of security from [7]. The solid black arrows denote implications whilst the dashed red arrows denote the opposite (i.e. separations).

1.3 Further Related Work

The investigation of relaxed, enhanced, and modified versions of CCA-2 security has a rich history and has found numerous applications in proxy-reencryption, updatable encryption, attribute based-encryption, rerandomizable encryption, or steganography [2,4,6,8,11,13,14,16,22,23].

The main relaxations of CCA-2, upon which the formalization of [7] builds, have been proposed in [24] as benign malleability and in [1] as generalized CCA-2 security, and also relate to loose ciphertext-unforgeability [17]. All these versions fall essentially into the formalization of public detectability discussed above, and all suffer from analogous technical issues, and hence in this work we focus on the formalization given in [7]. Three different flavours of RCCA have been introduced: IND-RCCA, UC-RCCA and NM-RCCA. In this work we focus on IND-RCCA. Our first benchmark is an abstract version of UC-RCCA. While the third flavour, NM-RCCA, is a strengthening of IND-RCCA (since it captures one additional attack vector), it does not seem to suffice to construct a confidential channel (or imply UC-RCCA for small message spaces) and is superseded in our treatment by IND-cl-RCCA that provably constructs the confidential channel for any message space.

A further relaxation of CCA-2 security, only loosely related to this work, is called detectable CCA-2 [15] and formalizes the detection of “dangerous” queries

in CCA-2 (without considering replayable properties). This notion provides a rather weak level of security on its own (in that it does not imply RCCA) [15].

Another line of research has consisted in studying *n-bounded* security definitions [10], i.e., a scheme may only be used to decrypt at most n messages. Crucially, this value n is known in advance and hardcoded in the scheme. Cramer et al. [10] showed that IND- n -bounded-CCA-2 is essentially equivalent to IND-CPA security.

2 Preliminaries

2.1 Constructive Cryptography

The Constructive Cryptography (CC) framework [19,18] is a composable security framework which views cryptography as a resource theory: a protocol transforms the assumed resources into the constructed resources.⁸ For example, if Alice and Bob have (access to) a shared secret key and an authentic channel, by running a one-time pad they construct a secure channel—this example is treated more formally further in this section.

In this view, encryption is the task of constructing channel resources. We thus start by defining various channels—used and constructed in this work—here below. Then we give the formal definition of a construction in CC.

INS. The weakest channel we consider is the (completely) insecure channel **INS**, where any message input by the sender goes straight to the adversary, and the adversary may insert any messages into the channel, which are then delivered to the receiver. This is drawn in the top left in Fig. 3.

AUT. In order to distribute the public keys used by PKE schemes, the players will also need an authentic channel **AUT**, which guarantees that anything received by the legitimate receiver was sent by the legitimate sender, but an adversary may also receive a copy of these messages. For simplicity, and without loss of generality, in our model we do not allow the adversary to either block an authentic channel or insert any replays. Such a channel is drawn in the top right of Fig. 3.

CONF. A confidential channel **CONF** only leaks the message length (denoted $|m|$) to the adversary, i.e. when the message m is input by the sender, the adversary receives $|m|$ at her interface. She can choose which message $j \leq i$ is delivered to the receiver, where i is the total number of messages input by the sender so far, or—since the channel is only confidential, but does not provide authenticity—the adversary may also inject a message of her own with (inj, m') , and m' is then delivered to the receiver. This is depicted in the bottom left of Fig. 3.

⁸ Resources essentially correspond to (ideal) functionalities in the Universal Composability framework (UC) [5], though in CC we additionally model the ability of players to communicate as having access to a channel resource.

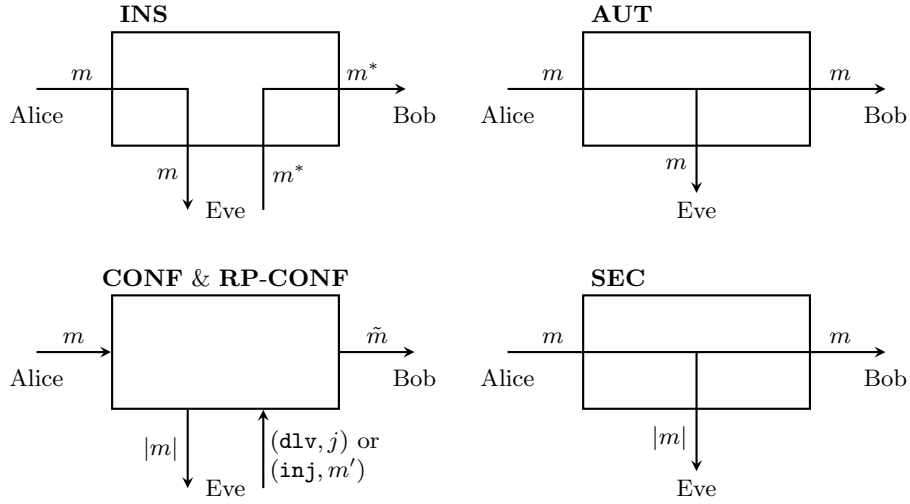


Fig. 3. A depiction of the channels used in this work. From top-left to bottom right: an insecure channel **INS**, an authentic channel **AUT**, a (replay protected) confidential channel (**RP**-)**CONF**, and a secure channel **SEC**.

RP-CONF. The **CONF** channel described above allows the adversary to deliver multiple times the same message to the receiver by inserting multiple times (dlv, j) . We define a stronger channel, the replay protected confidential channel **RP-CONF**, which will only process each (dlv, j) query at most once.

SEC. Finally, the secure channel **SEC** is both confidential and authentic, and is drawn in the bottom right of Fig. 3.

We will often consider channels that only transmit n messages, i.e., the sender may only input n messages and the adversary may only have n messages output to the receiver—either by delivering original messages or injecting new messages. These channels will be denoted $\mathbf{NAME}[n]$. The main properties of these channels are summarized in Fig. 4.

Formally, a *resource* (e.g. a channel) in an N -player setting is an interactive system with N interfaces, where each player may interact with the system at their interface by receiving outputs and providing inputs. These may be mathematically modeled as random systems [20,21] and can be specified by pseudo-code or an informal description as the channels above. In this work we consider the 3 player setting, and the interfaces are labeled A , B , and E for Alice, Bob, and Eve.

If multiple resources $\mathbf{R}_1, \dots, \mathbf{R}_\ell$ are accessible to players, we write $[\mathbf{R}_1, \dots, \mathbf{R}_\ell]$ the new resource resulting from having all resources accessible in parallel to the parties.

Operations run locally by some party (e.g. encrypting or decrypting a message) are modeled by interactive systems with two interface and are called *converters*.

Channel Name	Symbol	Leak $l(m)$	Insert	Replays
Insecure Channel	INS	m	Yes	Yes
Authentic Channel	AUT	m	No	No
Confidential Channel	CONF	$ m $	Yes	Yes
Replay Protected Confidential Channel	RP-CONF	$ m $	Yes	No
Secure Channel	SEC	$ m $	No	No

Fig. 4. A summary of the channel properties used in this work. *Leak* is the information about the message given to Eve, where $|m|$ denotes the length of the message. *Insert* denotes whether Eve is allowed to insert messages of her own into the channel. *Replay* denotes whether Eve can force a channel to deliver multiple times a message that was sent only once.

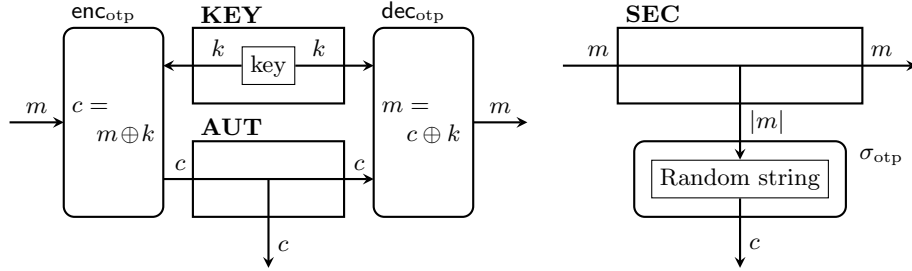


Fig. 5. The real and ideal systems for the one-time pad. Viewed as a black box, the real and ideal systems are indistinguishable.

The inner interface connects to the available resources, whereas the outer interface is accessible to the corresponding party to provide inputs and receive outputs. The composition of the resource and the converter is a new resource. For example, let \mathbf{R} be a resource, and let α be a converter which we connect at the A -interface of \mathbf{R} , then we write $\alpha^A \mathbf{R}$ for the new resource resulting from this connection. Formally, a converter is thus a map between resources.

To illustrate this, we draw the real system corresponding to a one-time pad encryption in Fig. 5. Here, the players have access to a secret key **KEY** and an authentic channel **AUT**. Alice runs the encryption converter enc_{otp} , which sends the ciphertext on the authentic channel. Bob runs the decryption converter dec_{otp} , which outputs the result of the decryption. The entire resource drawn on the left in Fig. 5 is denoted $\text{enc}_{\text{otp}}^A \text{dec}_{\text{otp}}^B [\mathbf{KEY}, \mathbf{AUT}]$, where the order of enc_{otp} and dec_{otp} does not matter since converters at different interfaces commute.

In order to argue that the protocol $\text{otp} = (\text{enc}_{\text{otp}}, \text{dec}_{\text{otp}})$ constructs a secure channel **SEC** from a shared secret key **KEY** and an authentic channel **AUT**, we need to find a converter σ_{otp} (called a *simulator*) such that when this simulator is attached to the adversarial interface of the constructed resource **SEC** (resulting in $\sigma_{\text{otp}}^E \mathbf{SEC}$), the real and ideal systems are indistinguishable. As illustrated in

Fig. 5, a simulator σ_{otp} which outputs a random string of the right length is sufficient for proving that the one-time pad constructs a secure channel.

Distinguishability between two systems \mathbf{R} and \mathbf{S} is defined with respect to a distinguisher \mathbf{D} which interacts with one of the systems, and has to output a bit corresponding to its guess. Let $\mathbf{D}[\mathbf{R}]$ and $\mathbf{D}[\mathbf{S}]$ denote the random variables corresponding to the output of \mathbf{D} when interacting with \mathbf{R} and \mathbf{S} , respectively. Then its advantage in distinguishing between the two is given by,

$$\Delta^{\mathbf{D}}(\mathbf{R}, \mathbf{S}) := |\Pr[\mathbf{D}[\mathbf{R}] = 0] - \Pr[\mathbf{D}[\mathbf{S}] = 0]|.$$

In the case of the one-time pad example with \mathbf{R} denoting the real system and \mathbf{S} the ideal system (drawn on the left and right in Fig. 5) we have that for all \mathbf{D} , $\Delta^{\mathbf{D}}(\mathbf{R}, \mathbf{S}) = 0$.

We now have all the elements needed to define a cryptographic construction in the three party setting.

Definition 1 (Asymptotic security [19,18]). *Let $\pi = \{(\pi_k^A, \pi_k^B)\}_{k \in \mathbb{N}}$ be an efficient family of converters, and let $\mathbf{R} = \{\mathbf{R}_k\}_{k \in \mathbb{N}}$ and $\mathbf{S} = \{\mathbf{S}_k\}_{k \in \mathbb{N}}$ be two efficient families of resources. We say that π asymptotically constructs \mathbf{R} from \mathbf{S} if there exists an efficient family of simulators $\sigma = \{\sigma_k\}_{k \in \mathbb{N}}$ such that for any efficient family of distinguishers $\mathbf{D} = \{\mathbf{D}_k\}_{k \in \mathbb{N}}$,*

$$\varepsilon(k) = \Delta^{\mathbf{D}_k}(\pi_k \mathbf{R}_k, \sigma_k \mathbf{S}_k)$$

is negligible. The construction is information-theoretically secure if the same holds for all (possibly inefficient) families of distinguishers.

For clarity we have made the security parameter k explicit in Definition 1, though in most of the technical part of this work we leave this parameter implicit to simplify the notation.

Remark 1. We note that defining the ideal world as the combination of simulator σ and an ideal resource \mathbf{S} , may be equivalently expressed as defining a set of acceptable ideal worlds $\{\sigma^E \mathbf{S}\}_\sigma$ over all possible σ . A construction is then secure if the real world (now also viewed as a set) is a subset of the ideal worlds. This may be generalized by considering different structures for the ideal world—not necessarily an ideal resource and a simulator. Though in this work we stick to the more traditional view of an ideal resource (or functionality) and simulator, since this generalization is not needed.

2.2 Public Key Encryption

The following definition of correctness for a PKE scheme is an adaptation of the almost-all-keys α -correctness definition from [12]. The main difference is that our definition requires for each single message in the message space correctness with high probability, whereas the notion from [12] only requires correctness with high probability over a uniform distribution on the message space. Thus, their notion allows for schemes which do not work correctly for certain plaintexts, while ours does not (and thus is suitable for composable frameworks).

Definition 2. A public key encryption (PKE) scheme Π with message space $\mathcal{M} \subseteq \{0, 1\}^*$, is a triple $\Pi = (G, E, D)$ of Probabilistic Polynomial-Time algorithms (PPTs) satisfying the following two conditions:

- On input 1^k (where k is the security parameter), G outputs a key-pair $(\mathbf{pk}, \mathbf{sk})$.
- With probability at least $1 - \text{negl}(k)$ (taken over the randomness of G), the key-pair $(\mathbf{pk}, \mathbf{sk})$ sampled by $G(1^k)$ according to the key-pair distribution is such that, for every $m \in \mathcal{M}$, $\Pr[D_{\mathbf{sk}}(E_{\mathbf{pk}}(m)) \neq m] \leq \text{negl}(k)$, where the probability is taken over the internal coin tosses of E and D .

Definition 3. An encryption scheme $\Pi = (G, E, D)$ is IND-CCA-2 secure if no PPT distinguisher \mathbf{D} distinguishes the two game systems $\mathbf{G}_0^{\Pi\text{-IND-CCA-2}}$ and $\mathbf{G}_1^{\Pi\text{-IND-CCA-2}}$ (specified below) with non-negligible advantage (in the security parameter k) over random guessing (i.e. if $\Delta^{\mathbf{D}}(\mathbf{G}_0^{\Pi\text{-IND-CCA-2}}, \mathbf{G}_1^{\Pi\text{-IND-CCA-2}}) \leq \text{negl}(k)$). For $b \in \{0, 1\}$, game system $\mathbf{G}_b^{\Pi\text{-IND-CCA-2}}$ is as follows:

Initialization: $\mathbf{G}_b^{\Pi\text{-IND-CCA-2}}$ generates a key-pair $(\mathbf{pk}, \mathbf{sk}) \leftarrow G(1^k)$, and sends \mathbf{pk} to \mathbf{D} .

First decryption stage: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system $\mathbf{G}_b^{\Pi\text{-IND-CCA-2}}$ computes $m = D_{\mathbf{sk}}(c)$ and sends m to \mathbf{D} .

Challenge stage: When \mathbf{D} queries $(\text{test messages}, m_0, m_1)$, for $m_0, m_1 \in \mathcal{M}$ such that $|m_0| = |m_1|$, $\mathbf{G}_b^{\Pi\text{-IND-CCA-2}}$ computes $c^* = E_{\mathbf{pk}}(m_b)$, and sends c^* to \mathbf{D} .⁹

Second decryption stage: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system $\mathbf{G}_b^{\Pi\text{-IND-CCA-2}}$ replies **test** if $c = c^*$ and replies $m = D_{\mathbf{sk}}(c)$ (i.e. the decryption of c) otherwise.

For simplicity, throughout the paper we will omit the prefix Π from the notation of the game systems, unless needed for clarity.

Definition 4. An encryption scheme $\Pi = (G, E, D)$ is IND-RCCA secure if it is secure according to the definition of IND-CCA-2 security (Definition 3), but where the IND-RCCA game systems differ from the IND-CCA-2 game systems in the second decryption stage, which now works as follows: In the following, let m_0, m_1 be the two challenge messages queried by distinguisher \mathbf{D} during the Challenge stage:

Second decryption stage: When \mathbf{D} queries $(\text{ciphertext}, c)$, the game system computes $m = D_{\mathbf{sk}}(c)$. If $m \in \{m_0, m_1\}$, then the game system replies with the special response **test** to \mathbf{D} , and otherwise sends m to \mathbf{D} .

2.3 Public Key Encryption With Replay Filtering

We now introduce two new types of PKE schemes, namely ones in which ciphertext replays can be efficiently detected by an algorithm F that is defined as part of the scheme. For the correctness condition of these schemes we require, in addition

⁹ Unless explicitly stated, we assume that \mathbf{D} can only perform a single challenge query.

to the usual correctness condition of PKE schemes, that with high probability F cannot relate two fresh encryptions of any messages. This is an essential requirement such that F can be used for filtering out ciphertext replays, because the correctness condition guarantees that it will not filter out honestly generated ciphertexts (later in Section 5.2 we couple such schemes with the proper security notions).

Definition 5. A PKE scheme with secret replay filtering (PKESF) Π with message space $\mathcal{M} \subseteq \{0, 1\}^*$, is a 4-tuple $\Pi = (G, E, D, F)$ of Probabilistic Polynomial-Time algorithms (PPTs) satisfying the following two conditions:

- On input 1^k (where k is the security parameter), G outputs a key-pair $(\mathbf{pk}, \mathbf{sk})$.
- With probability at least $1 - \text{negl}(k)$ (taken over the randomness of G), the key-pair $(\mathbf{pk}, \mathbf{sk})$ sampled by $G(1^k)$ according to the key-pair distribution is such that, $\forall m, m' \in \mathcal{M}$,

$$\Pr[D_{\mathbf{sk}}(E_{\mathbf{pk}}(m)) \neq m \vee F(\mathbf{pk}, \mathbf{sk}, E_{\mathbf{pk}}(m), E_{\mathbf{pk}}(m')) = 1] \leq \text{negl}(k),$$

where the probability is taken over the internal coin tosses of E , D and F .

A public key encryption scheme with public replay filtering (PKEPF) Π is just like a PKESF except that F now does not receive the secret key \mathbf{sk} .

As one might note, from any correct and IND-CCA-2 secure PKE scheme $\Pi = (G, E, D)$, one can define a correct PKEPF scheme $\Pi' = (G, E, D, F)$ where $F(\mathbf{pk}, c, c') = 1$ if and only if $c = c'$; the correctness of Π' with respect to Definition 5 follows from the correctness and IND-CCA-2 security of Π .

2.4 Reductions

Most of the proofs in this work consist in showing reductions between various security definitions. Both the constructive statements introduced in Sect. 2.1 and game-based definitions such as IND-CCA-2 (Definition 3) can be viewed as distinguishing systems—the real world \mathbf{W}_0 from the ideal world \mathbf{W}_1 and game \mathbf{G}_0 from game \mathbf{G}_1 , respectively. A reduction between two such definitions consists in proving that if a distinguisher \mathbf{D} can succeed in one task, then a (related) distinguisher \mathbf{D}' can succeed in the other. We only give explicit reductions with single blackbox access to \mathbf{D} in this work, i.e., we define $\mathbf{D}' := \mathbf{DC}$, where \mathbf{DC} denotes the composition of two systems \mathbf{D} and \mathbf{C} . \mathbf{C} is called the *reduction system* (or simply the *reduction*).

For example, if we wish to reduce the task of breaking a constructive definition (with real and ideal systems $\mathbf{W}_0 = \pi^{AB}\mathbf{R}$ and $\mathbf{W}_1 = \sigma^E\mathbf{S}$ for some simulator σ) to a game-based definition (with games \mathbf{G}_0 and \mathbf{G}_1), we will typically fix σ and find a system \mathbf{C} such that $\mathbf{W}_0 = \mathbf{CG}_0$ and $\mathbf{W}_1 = \mathbf{CG}_1$. Then

$$\Delta^{\mathbf{D}}(\mathbf{W}_0, \mathbf{W}_1) = \Delta^{\mathbf{D}}(\mathbf{CG}_0, \mathbf{CG}_1) = \Delta^{\mathbf{DC}}(\mathbf{G}_0, \mathbf{G}_1),$$

i.e., given a distinguisher \mathbf{D} that can distinguish \mathbf{W}_0 from \mathbf{W}_1 with non-negligible advantage, we get an explicit new distinguisher \mathbf{DC} that can win the game with

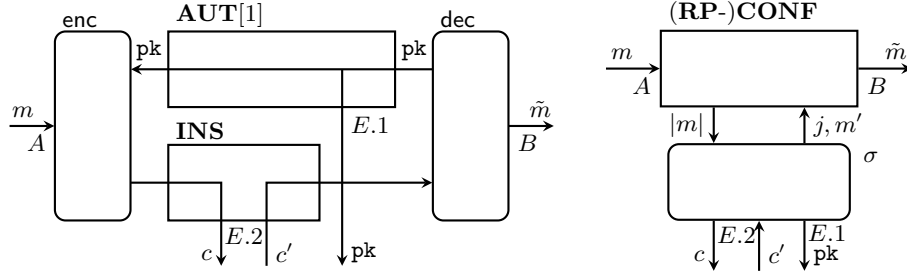


Fig. 6. Real and ideal systems for (replay protected) confidential channel construction. Capital letters ($A, B, E.1, E.2$) represent interface labels and small letters ($m, \tilde{m}, c, c', j, \text{pk}$) represent values that are in- or output.

non-negligible advantage. Or, the contrapositive, if \mathbf{G}_0 and \mathbf{G}_1 are hard to distinguish, then in particular they are hard to distinguish for all distinguishers of the form \mathbf{DC} (for any efficient \mathbf{D} and fixed \mathbf{C}). This means that no efficient distinguisher \mathbf{D} can tell \mathbf{W}_0 from \mathbf{W}_1 for the given simulator σ .

3 Benchmarking Confidentiality

In this section we present three benchmark constructions to capture the security of confidential communication and replay protected confidential communication.

3.1 Benchmark 1: The CONF Channel

The first channel we want to construct is the confidential channel **CONF** introduced in Sect. 2.1. The ideal system thus simply consists of this channel and a simulator σ , as depicted on the right in Fig. 6, and is denoted $\sigma^E \mathbf{CONF}$.

In order to achieve this, Alice and Bob need an authentic channel for one message **AUT[1]** (from Bob to Alice), so that Bob can send his public key authentically to Alice. They also use a completely insecure channel **INS** to transmit the ciphertexts. Alice's converter **enc** encrypts any messages with the public key obtained from **AUT[1]**, and sends the resulting ciphertext on **INS** (i.e., for a PKE $\Pi = (G, E, D)$, **enc** runs E). Bob's converter **dec** generates the key-pair (pk, sk) , sends **pk** over **AUT[1]** to Alice, and decrypts any ciphertext received from **INS** using **sk** (i.e., **dec** runs G and D). The resulting message is output at Bob's outer interface B (to the environment/distinguisher). This real system is drawn on the left in Fig. 6), and is denoted $\text{enc}^A \text{dec}^B [\mathbf{AUT}[1], \mathbf{INS}]$.

In the case where the key-pair is used at most n times, the channels used and constructed are limited to n uses, which we then denote **INS** $[n]$ and **CONF** $[n]$, respectively.

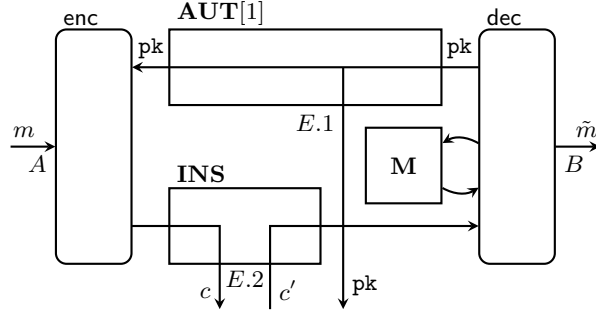


Fig. 7. Real system for constructing a replay protected confidential channel. Capital letters (A , B , $E.1$, $E.2$) represent interface labels and small letters (m , \hat{m} , c , c' , pk) represent values that are in- or output.

3.2 Benchmark 2: The RP-CONF Channel

As explained in Sect. 1.2, our second benchmark is the construction of a stronger channel, namely a replay protected confidential channel, i.e. one in which an adversary's input (dlv, j) may only be processed once for each j . The ideal system $\sigma^E \mathbf{RP-CONF}$ is thus similar to the one of Benchmark 1, only differing in the underlying ideal channel which now is the stronger **RP-CONF** channel.

The real system is similar to the real system from Benchmark 1 in that we want to construct **RP-CONF** from a single use authentic channel $\mathbf{AUT}[1]$ and an insecure channel **INS**. However, the replay detection algorithm requires memory to store the ciphertexts it has already processed. We model this memory use explicitly by providing a memory resource \mathbf{M} to the decryption converter. This is drawn in Fig. 7. The real system is thus $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}, \mathbf{M}]$.

If one uses a public key encryption scheme with replay filtering defined by an algorithm F (see Sect. 2.3), then Alice's converter enc runs the encryption algorithm as for a normal PKE, but Bob's converter additionally runs the filtering algorithm F before decrypting to detect replays.

We will often parameterize the channels by the number of messages n input at Alice's interface, which we denote $\mathbf{INS}[n]$ and $\mathbf{RP-CONF}[n]$, respectfully. The distinguishability between real and ideal systems is then expressed as a function of n . The encryption and decryption functions considered are independent of this value, and thus by taking the limit for $n \rightarrow \infty$ one gets asymptotic security bounds for a given scheme. We point out that this stands in sharp contrast with the n -bounded security notions from [10], which are much weaker, since the value n (of the allowed *decryption* queries) is known in advance and hardcoded in the scheme.

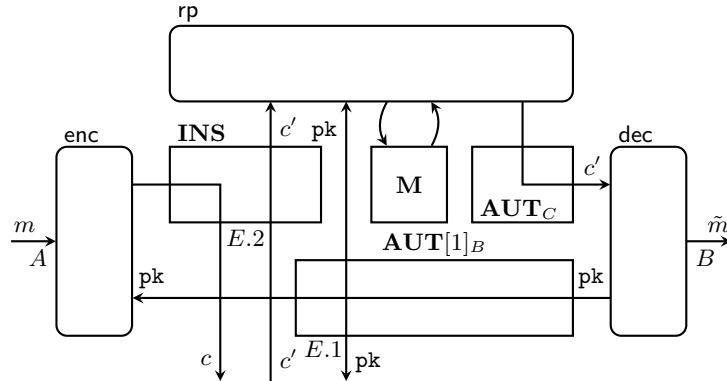


Fig. 8. Real system for constructing a reply protected confidential channel with outsourced replay filtering. As in previous figures, the sender Alice is on the left, the receiver Bob is on the right and the eavesdropper Eve is below. In this setting we have another party, Charlie, above in the picture, to whom replay detection is outsourced, and who runs the converter rp . Capital letters ($A, B, E.1, E.2$) represent interface labels and small letters (m, \tilde{m}, c, c', pk) represent values that are in- or output.

3.3 Benchmark 3: The RP-CONF Channel With Outsourcable Replay Protection

In this section we again want to construct a replay protected confidential channel **RP-CONF**—but where the job of filtering out ciphertext replays is outsourced to a third party. The ideal system is thus identical to Benchmark 2, i.e $\sigma^E \mathbf{RP-CONF}$.

The real system now has three honest parties, Alice the sender, Bob the receiver, and Charlie the replay-filterer, where each runs its own converter enc , dec and rp , respectively. As before, a public key pk is generated by dec and sent on an authentic channel $AUT[1]_B$ to both Alice and Charlie—but Eve gets a copy as well—where the index B denotes the origin of the authenticated message. And as before, enc encrypts the message and sends it on an insecure channel INS , but this time Charlie is on the receiving end of INS . Charlie then runs rp , which decides if the message should be forwarded to Bob on a channel AUT_C or if it gets deleted—this channel needs to be authenticated, so that Eve cannot change the messages or inject replays again. To do this, rp needs access to the memory resource M , which stores the previous ciphertexts. Finally, dec decrypts the messages received. This is depicted in Fig. 8.

Note that in this setup, rp does not have access to the secret key and so it must detect replays with the public key only; and dec does not have access to the memory M , and can thus not perform the replay filtering itself. In the case where the players use a PKEPF $\Pi = (G, E, D, F)$, then enc runs E , dec runs G and D , and rp runs F .

As before, we will often parameterize the channels by the number of messages n input at Alice’s interface, which is independent of the encryption and decryption functions.

4 Technical Issues with pd-RCCA and sd-RCCA

In [7], Canetti et al. introduce pd-RCCA and sd-RCCA as supposedly relaxed versions of CCA-2 security. Although other supposedly relaxed versions of CCA-2, such as Benign Malleability [24] and generalized CCA-2 security [1], had been introduced before, these notions are subsumed by the definition of pd-RCCA and suffer from the same technical issues we uncover in this section. For this reason, we will focus only on the pd-RCCA and sd-RCCA security notions. We now recall the definition of IND-pd-RCCA and IND-sd-RCCA [7].

Definition 6. Let $\Pi = (G, E, D)$ be an encryption scheme.

1. Say that a family of binary relations $\equiv_{\mathbf{pk}}$ (indexed by the public keys of Π) on ciphertext pairs is a compatible relation for Π if for all key-pairs $(\mathbf{pk}, \mathbf{sk})$ of Π :
 - (a) For any two ciphertexts c, c' , if $c \equiv_{\mathbf{pk}} c'$, then $D_{\mathbf{sk}}(c) = D_{\mathbf{sk}}(c')$, except with negligible probability over the random choices of D .
 - (b) For any plaintext $m \in \mathcal{M}$, if c and c' are two ciphertexts obtained as independent encryptions of m (i.e. two applications of algorithm E on m using independent random bits), then $c \equiv_{\mathbf{pk}} c'$ only with negligible probability.
2. We say that a relation family as above is publicly computable (resp. secretly computable) if for all key pairs $(\mathbf{pk}, \mathbf{sk})$ and ciphertext pairs (c, c') it can be determined whether $c \equiv_{\mathbf{pk}} c'$ using a PPT algorithm taking inputs (\mathbf{pk}, c, c') (resp. $(\mathbf{pk}, \mathbf{sk}, c, c')$).
3. We say that Π is publicly-detectable Replayable-CCA (IND-pd-RCCA) if there exists a compatible and publicly computable relation family $\equiv_{\mathbf{pk}}$ such that Π is secure according to the standard definition of IND-CCA-2 (Definition 3), but where the game systems differ from the IND-CCA-2 game systems in the second decryption stage, which now works as follows: In the following, let c^* be the challenge ciphertext output by the game system:

Second decryption stage: When \mathbf{D} queries $(\text{ciphertext}, c)$, the game system replies **test** if $c^* \equiv_{\mathbf{pk}} c$, and otherwise computes $m = D_{\mathbf{sk}}(c)$ and then sends m to \mathbf{D} .

Similarly, we say that Π is secretly-detectable Replayable-CCA (IND-sd-RCCA) if the above holds for a secretly computable relation family $\equiv_{\mathbf{pk}}$.

It is claimed in [7] that IND-CCA-2 security implies IND-pd-RCCA security (with the equality relation serves as the compatible relation), which in turn implies IND-sd-RCCA security. However, as we now show, Definition 6 is not an actual relaxation of the IND-CCA-2 security notion. More concretely, we prove that IND-CCA-2 security does not entail IND-pd-RCCA nor even IND-sd-RCCA security, according to their definition.

Theorem 1. *If there is a correct and IND-CCA-2 secure PKE scheme, then there is a correct and IND-CCA-2 secure PKE scheme which is not IND-pd-RCCA nor IND-sd-RCCA secure.*

Throughout the rest of the section, let $\Pi = (G, E, D)$ be a correct and IND-CCA-2 secure PKE scheme. Without loss of generality, assume that all messages in Π 's message space have the same length. We create a scheme $\Pi' = (G', E', D')$ (see Algorithm 1) such that Π' is a correct and IND-CCA-2 secure PKE scheme, but is not IND-pd-RCCA nor IND-sd-RCCA secure.

Algorithm 1 The Π' scheme.

1: procedure $G'(1^n)$ 2: $(\text{pk}, \text{sk}) \leftarrow G(1^n)$ 3: $\tilde{m} \leftarrow^{\$} \mathcal{M}$ 4: $c \leftarrow E_{\text{pk}}(\tilde{m})$ 5: return $(\text{pk}', \text{sk}') \leftarrow (\text{pk}, (\text{sk}, c))$ 6: end procedure 7: procedure $E'_{\text{pk}':=\text{pk}}(m)$ 8: return $E_{\text{pk}}(m)$ 9: end procedure	10: procedure $D'_{\text{sk}':=(\text{sk},c)}(c)$ 11: if $\text{sk}'.c \neq c$ then 12: return $D_{\text{sk}}(c)$ 13: else 14: $b \leftarrow^{\$} \{0, 1\}$ 15: if $b = 0$ then 16: return \perp 17: else 18: return $D_{\text{sk}}(c)$ 19: end if 20: end if 21: end procedure
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Lemma 1. *If Π is correct and IND-CCA-2 secure, then so is Π' .*

Proof. It is easy to see that if Π is correct and IND-CCA-2 secure then Π' is a correct PKE scheme. We now prove that Π' is IND-CCA-2 secure.

Let \mathbf{D} be a distinguisher for the IND-CCA-2 game systems for Π' . We construct a distinguisher \mathbf{D}' , which internally uses \mathbf{D} , for the IND-CCA-2 game systems for Π such that

$$\Delta^{\mathbf{D}'}(\mathbf{G}_0^{\Pi\text{-IND-CCA-2}}, \mathbf{G}_1^{\Pi\text{-IND-CCA-2}}) = \Delta^{\mathbf{D}}(\mathbf{G}_0^{\Pi'\text{-IND-CCA-2}}, \mathbf{G}_1^{\Pi'\text{-IND-CCA-2}}). \quad (4.1)$$

\mathbf{D}' works as follows: When \mathbf{D}' receives pk from the game, it picks a plaintext \tilde{m} uniformly at random from \mathcal{M} , generates a ciphertext $c = E_{\text{pk}}(\tilde{m})$, and forwards pk to \mathbf{D} . Before the challenge ciphertext is set, whenever \mathbf{D} queries $(\text{ciphertext}, c')$, \mathbf{D}' first checks if $c = c'$: if this is the case then \mathbf{D}' flips a coin uniformly at random and (depending on the outcome of the coin) either returns \perp as the result of the query, or forwards it to the IND-CCA-2 game. If $c \neq c'$ then \mathbf{D}' simply forwards the query to the game. Upon receiving the result of the decryption query, \mathbf{D}' forwards it to \mathbf{D} . When \mathbf{D} issues the challenge query, \mathbf{D}' forwards it to the game, and, upon receiving the challenge ciphertext c^* from the game, \mathbf{D}' forwards it back to \mathbf{D} . After the challenge ciphertext is set, whenever \mathbf{D} issues a

decryption query $(\text{ciphertext}, c')$, \mathbf{D}' behaves just as before, unless $c' = c^*$. In such case, \mathbf{D}' simply forwards the decryption query to the IND-CCA-2 game and returns the result to \mathbf{D} . When \mathbf{D} outputs a guess b , \mathbf{D}' outputs the same guess and terminates. Clearly, 4.1 holds, and thus, if Π is IND-CCA-2 secure, then so is Π' . \square

We now show that a compatible relation for Π' cannot relate any freshly generated ciphertext to itself.

Lemma 2. *Let \equiv_{pk} be any family of compatible relations for Π' (indexed by the public keys of Π'). Then, for each pk in the support of Π' 's public keys, we have: for any fresh encryption c of some plaintext $m \in \mathcal{M}$ under pk , $c \not\equiv_{\text{pk}} c$.*

Proof. For each public key pk in the support of Π' 's public keys, let \equiv_{pk} be a compatible relation of Π' with respect to pk . For each ciphertext c that can be generated as a fresh encryption of some plaintext m by E' under pk , there is a key-pair (pk, sk) (for the same public key pk) such that $\Pr[D'_{\text{sk}}(c) \neq D_{\text{sk}}(c)] \geq \frac{1}{2}$. Hence, by the compatibility condition of Definition 6, $c \not\equiv_{\text{pk}} c$. \square

Lemma 3. *Π' is not IND-pd-RCCA nor IND-sd-RCCA secure.*

Proof. By the definitions of IND-pd-RCCA and IND-sd-RCCA, the challenge ciphertext c^* is always a fresh encryption of some plaintext. By Lemma 2 it then follows $c^* \not\equiv_{\text{pk}} c^*$. As such, a distinguisher is allowed to simply ask for the decryption of the challenge c^* and thus distinguish the two game systems. \square

Lemmas 1 and 3 conclude the proof of Theorem 1.

One way to solve this technical issue with the definitions of IND-pd-RCCA and IND-sd-RCCA is by restricting the class of schemes we consider. For instance, if we would require the decryption algorithm of PKE schemes to be deterministic, then our counterexample would not apply, and the IND-pd-RCCA and IND-sd-RCCA security notions would indeed be implied by IND-CCA-2 security. However, it still seems that there should be a concept of ciphertext relation that is conceptually sound for any encryption scheme.

5 Relaxing Chosen Ciphertext Security

As discussed in Sect. 1, while IND-CCA-2 is generally a too strong security notion, IND-RCCA security is too weak, in that it is not sufficient to achieve the weaker Benchmark 1 for small message spaces. In this section we introduce three new security notions—which are provably between IND-CCA-2 and IND-RCCA, see Sect. 6—and prove that they are sufficient to achieve the three benchmarks introduced in Sect. 3.

5.1 Achieving Benchmark 1: Constructing the CONF Channel

A game-based security notion that captures the confidentiality of an encryption scheme against active adversaries is one which is sufficiently strong to achieve a confidential channel (as defined in Sect. 3.1). Yet, it must also be as weak as possible so that it does not exclude any schemes which provide confidentiality. To achieve this, we introduce the IND-cl-RCCA security notion, and its multi-challenge version $[n]$ IND-cl-RCCA.

Definition 7. We say that an encryption scheme $\Pi = (G, E, D)$ is IND-cl-RCCA secure if there exists an efficient algorithm v that takes as input a key-pair (pk, sk) and a pair of ciphertexts c, c' and outputs a boolean (corresponding to whether the ciphertexts seem related or not), such that no PPT distinguisher \mathbf{D} distinguishes the game systems $\mathbf{G}_0^{\text{IND-cl-RCCA}}$ and $\mathbf{G}_1^{\text{IND-cl-RCCA}}$ (specified below) with non-negligible advantage (in the security parameter k) over random guessing. For $b \in \{0, 1\}$, game system $\mathbf{G}_b^{\text{IND-cl-RCCA}}$ is as follows:

Initialization: $\mathbf{G}_b^{\text{IND-cl-RCCA}}$ generates a key-pair $(\text{pk}, \text{sk}) \leftarrow G(1^k)$, and sends pk to \mathbf{D} .

First decryption stage: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system $\mathbf{G}_b^{\text{IND-cl-RCCA}}$ computes $m = D_{\text{sk}}(c)$ and sends m to \mathbf{D} .

Challenge stage: When \mathbf{D} queries $(\text{test messages}, m_0, m_1)$, for $m_0, m_1 \in \mathcal{M}$ such that $|m_0| = |m_1|$, $\mathbf{G}_b^{\text{IND-cl-RCCA}}$ computes $c^* = E_{\text{pk}}(m_b)$, and sends c^* to \mathbf{D} .

Second decryption stage: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system $\mathbf{G}_b^{\text{IND-cl-RCCA}}$ calls $v(\text{pk}, \text{sk}, c^*, c)$ and decrypts c , obtaining a plaintext $m = D_{\text{sk}}(c)$. If v 's output is 1 and $m = m_b$, the game system replies **test** to \mathbf{D} , and in all other cases the game replies with m .

At a high level, the job of algorithm v is to disallow strategies that an adversary could take to win the security game, but would not help break confidentiality of the encryption. In the context of the IND-cl-RCCA game, v is used to disallow adversaries to pursue strategies in which they would ask for the decryption of a ciphertext that would decrypt to the challenge message (a so-called *replay*). Thus, the game can only refuse to answer a decryption query for a ciphertext c if both of the following two conditions are met: 1. according to v , c is a replay of the challenge ciphertext; and 2. c indeed decrypts to the same plaintext as the challenge ciphertext.

IND-cl-RCCA security is sufficient for achieving Benchmark 1 for a single message (i.e. constructing an ideal **CONF**[1] channel)—this follows from Theorem 2 below. However, it is not clear whether it is also sufficient for achieving Benchmark 1 for multiple messages: since, in order to check if two ciphertexts are related, v requires the secret key, it becomes apparently unfeasible to detect relations between pairs of arbitrary ciphertexts, which is crucial for making a hybrid reduction from distinguishing $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]]$ from **CONF**[n] to distinguishing the two IND-cl-RCCA game systems. To achieve Benchmark 1 for multiple messages, we now present the multi-challenge version of IND-cl-RCCA

security, which we denote by $[n]$ IND-cl-RCCA security, where n is the maximum number of challenge queries that a distinguisher can make.

Definition 8. We say that an encryption scheme $\Pi = (G, E, D)$ is $[n]$ IND-cl-RCCA secure if it is secure according to Definition 7, but where, for $b \in \{0, 1\}$, the game system $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$, which now accepts n challenge queries, behaves as follows:

Initialization: First, $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ creates and initializes a table t of plaintext-ciphertext pairs which is initially empty. Then, $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ runs $(\mathbf{pk}, \mathbf{sk}) \leftarrow G(1^k)$, and sends \mathbf{pk} to \mathbf{D} .

Decryption queries: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system calls, for each plaintext-ciphertext pair $(m_{b,j}, c_j^*)$ stored in t , $v(\mathbf{pk}, \mathbf{sk}, c_j^*, c)$ and decrypts c , obtaining a plaintext $m = D_{\mathbf{sk}}(c)$. If for every plaintext-ciphertext pair stored in t , either v 's output is 0 or $m \neq m_{b,j}$, then the game system replies with m to \mathbf{D} . Otherwise, let $(m_{b,l}, c_l^*)$ be the plaintext-ciphertext pair stored in t with the smallest l such that both $v(\mathbf{pk}, \mathbf{sk}, c_l^*, c) = 1$ and $m = m_{b,l}$. Then, $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ replies (test, l) to \mathbf{D} .

i -th challenge query (for $i \leq n$): Whenever the distinguisher \mathbf{D} issues a challenge query $(\text{test messages}, m_{0,i}, m_{1,i})$, where $m_{0,i}, m_{1,i} \in \mathcal{M}$ such that $|m_{0,i}| = |m_{1,i}|$, the game system computes $c_i^* = E_{\mathbf{pk}}(m_{b,i})$, stores $(m_{b,i}, c_i^*)$ in table t , and sends c_i^* to \mathbf{D} .

We now show that $[n]$ IND-cl-RCCA security is sufficient for achieving Benchmark 1 for n messages. Thus, we need to prove that the construction is indistinguishable from the ideal $\mathbf{CONF}[n]$ channel up to the $[n]$ IND-cl-RCCA security of the underlying PKE scheme.

Let $\Pi = (G, E, D)$ be a correct and $[n]$ IND-cl-RCCA secure PKE scheme, and let the protocol $\pi = (\text{enc}, \text{dec})$ be such that Alice's converter enc runs the encryption algorithm E at most n times, and Bob's converter dec generates the initial pair of secret and public keys with G , sends \mathbf{pk} on the authentic channel to Alice and runs D on the ciphertexts received at most n times.

To prove that π constructs $\mathbf{CONF}[n]$ from $\mathbf{AUT}[1]$ and $\mathbf{INS}[n]$ (Definition 1), we show how to create, from any algorithm v that satisfies Definition 8, an efficient simulator σ which internally uses v such that any distinguisher \mathbf{D} for $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]]$ and $\sigma^E \mathbf{CONF}[n]$ can be transformed into an equally good distinguisher for the $[n]$ IND-cl-RCCA game systems. Then, from the $[n]$ IND-cl-RCCA security of Π , it follows that there is such an algorithm v , implying that no efficient distinguisher \mathbf{D} can distinguish between the real world $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]]$ and the ideal world $\sigma^E \mathbf{CONF}[n]$ with simulator σ attached. In turn, this implies that Benchmark 1 is achieved.

Theorem 2. Let v be an algorithm that suits $[n]$ IND-cl-RCCA (Definition 8). There exists an efficient simulator σ and an efficient reduction \mathbf{R} such that for

every distinguisher \mathbf{D} ,

$$\begin{aligned} \Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]], \sigma^E \mathbf{CONF}[n]) \\ = \Delta^{\mathbf{DR}}(\mathbf{G}_0^{[n]\text{IND-cl-RCCA}}, \mathbf{G}_1^{[n]\text{IND-cl-RCCA}}). \end{aligned}$$

Proof. Consider the following simulator σ for interface E of $\mathbf{CONF}[n]$, which has two sub-interfaces denoted by $E.1$ and $E.2$ on the outside (since the real-world system also has two sub-interfaces at E): Initially, σ generates a key-pair $(\mathbf{pk}, \mathbf{sk})$ and outputs \mathbf{pk} at $E.1$. When it receives the i -th input l_i at the inside interface in (which is connected to $\mathbf{CONF}[n]$), σ generates an encryption $c \leftarrow E_{\mathbf{pk}}(\tilde{m})$ of a randomly chosen message \tilde{m} of length l , records (i, \tilde{m}, c) and outputs c at $E.2$. When c' is input at $E.2$, σ proceeds as follows: First, it decrypts c' , obtaining some plaintext m' . If (j, \tilde{m}, c) has been recorded for some j such that $\tilde{m} = m'$ and $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$, then σ outputs (dlv, j) at in (where j is the smallest index satisfying this condition). If no such triple has been recorded, σ outputs (inj, m') at in (unless $m' = \perp$).

Having defined the simulator σ , we now introduce a reduction system \mathbf{R} , such that for any efficient distinguisher \mathbf{D}

1. $\mathbf{RG}_0^{[n]\text{IND-cl-RCCA}} \equiv \text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]]$; and
2. $\mathbf{RG}_1^{[n]\text{IND-cl-RCCA}} \equiv \sigma^E \mathbf{CONF}[n]$.

Consider the following reduction system \mathbf{R} (which processes at most n inputs at the outside A interface): Initially, \mathbf{R} forwards the public key \mathbf{pk} generated by the game system to the $E.1$ interface. When the j -th message m is input at the A interface of \mathbf{R} : \mathbf{R} chooses a message \tilde{m} of length $|m|$ uniformly at random, and makes the challenge query $(\text{test messages}, m, \tilde{m})$ to the game system, which replies with some ciphertext c . Then, \mathbf{R} records $m_j^* = m$. Next, \mathbf{R} outputs c at the outside $E.2$ interface. When (inj, c') is input at interface $E.2$, \mathbf{R} behaves as follows. First, \mathbf{R} makes a decryption query for c' to the game, obtaining some m' . If $m' = (\text{test}, j)$, then \mathbf{R} outputs m_j^* at interface B . If $m' = \perp$, \mathbf{R} ignores the injection, and nothing happens. Else, \mathbf{R} outputs m' at the B interface. It is easy to see that indeed $\mathbf{RG}_0^{[n]\text{IND-cl-RCCA}} \equiv \text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]]$ and $\mathbf{RG}_1^{[n]\text{IND-cl-RCCA}} \equiv \sigma^E \mathbf{CONF}[n]$. Using the above facts, it finally follows

$$\begin{aligned} \Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n]], \sigma^E \mathbf{CONF}[n]) \\ = \Delta^{\mathbf{D}}(\mathbf{RG}_0^{[n]\text{IND-cl-RCCA}}, \mathbf{RG}_1^{[n]\text{IND-cl-RCCA}}) \\ = \Delta^{\mathbf{DR}}(\mathbf{G}_0^{[n]\text{IND-cl-RCCA}}, \mathbf{G}_1^{[n]\text{IND-cl-RCCA}}). \quad \square \end{aligned}$$

5.2 Achieving Benchmark 2: Constructing the RP-CONF Channel

Another use of IND-CCA-2 security is for achieving replay protected confidential communication. As hinted by Benchmarks 2 and 3, replay protection comes in two flavours: 1. private detection and filtering of replays; and 2. public detection

and filtering of replays. We begin by looking into the setting where Bob is the one responsible for filtering out ciphertext replays (Benchmark 2).

Before introducing a new security notion, we first look into why IND-cl-RCCA does not seem to suffice for constructing the **RP-CONF** channel. First, note that, the **RP-CONF** channel construction (Benchmark 2) has to protect not only against replays of ciphertexts sent by Alice, but also against replays of ciphertexts injected by Eve. This is so since the receiving end (i.e. the dec converter) does not know where the ciphertexts have originated.¹⁰ Hence, for each ciphertext that the converter receives, it has to make sure that it is not a replay of any previously received ciphertext, implying that the converter has to impede all ciphertext replays. When one tries to make a reduction from distinguishing the real world construction $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}, \mathbf{M}]$ and the ideal world channel **RP-CONF** to winning the IND-cl-RCCA game, two critical issues arise:

1. The algorithm v used by the game systems might not compute an equivalence relation: Consider the case where Alice inputs a message m into the channel which results in a ciphertext c being output at the E interface. Eve can create two distinct replays of the ciphertext c , say c' and c'' , and input them into the E interface. While, from IND-cl-RCCA security, v should detect that ciphertext c is related to both c' and c'' , it does not necessarily detect whether c' is related to c'' . In such case, v cannot be used to detect ciphertext replays, as it would allow Eve to replay what Alice sends, by generating different replays of c and injecting them into the channel (without ever injecting c into the channel).
2. The reduction does not have access to the secret key generated by the game system: Even assuming that v computes an equivalence relation, it is not clear how one could reduce distinguishing the real and ideal worlds to distinguishing the two underlying IND-cl-RCCA game systems. Since any reduction system \mathbf{R} that one would attach to the game systems does not have access to the secret key, it is not clear how \mathbf{R} would be able to check if any arbitrary pair of ciphertexts c' and c'' are related according to v (i.e. \mathbf{R} would be able to compute $v(\mathbf{pk}, \mathbf{sk}, c', c'')$ without knowing \mathbf{sk}).

Interestingly these remarks also apply to the IND-sd-RCCA notion from [7], hinting at the fact that the IND-sd-RCCA security notion does not capture what it was meant to capture. Another interesting remark is that, as for IND-cl-RCCA, the single challenge and the multi challenge versions of IND-sd-RCCA security do not seem to be necessarily equivalent.¹¹ With this, we now introduce IND-srp-RCCA security, which captures the secret detectability of ciphertext replays.

Definition 9. *A scheme $\Pi = (G, E, D)$ is IND-srp-RCCA secure if there exists an efficient algorithm v that computes, for each key-pair $(\mathbf{pk}, \mathbf{sk})$, an equivalence relation over ciphertexts c, c' such that for every key-pair $(\mathbf{pk}, \mathbf{sk})$ in the*

¹⁰ Note that, other than the assumption that the public key is authentically transmitted, we are only assuming an insecure channel between Alice and Bob.

¹¹ We leave the problem of proving whether these notions are equivalent or not as open.

support of $G(1^k)$ and every pair of ciphertexts c, c' , if $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$ then $\delta(D_{\mathbf{sk}}(c), D_{\mathbf{sk}}(c')) \leq \text{negl}(k)$ (where the randomness is over the internal randomness of D), and if no efficient distinguisher \mathbf{D} distinguishes the game systems $\mathbf{G}_0^{\text{IND-srp-RCCA}}$ and $\mathbf{G}_1^{\text{IND-srp-RCCA}}$ (specified below) with non-negligible advantage (in the security parameter k) over random guessing. The IND-srp-RCCA game systems work just as the IND-CCA-2 game systems, except that the IND-srp-RCCA game systems give distinguisher \mathbf{D} oracle access to v throughout the entire game (so that \mathbf{D} can check whether any two ciphertexts c, c' are related according to v with respect to the key-pair \mathbf{pk}, \mathbf{sk} generated by the game system), and also except for the second decryption stage, which now works as follows:

Second decryption stage: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system replies `test` if $v(\mathbf{pk}, \mathbf{sk}, c^*, c) = 1$ and replies $m = D_{\mathbf{sk}}(c)$ otherwise.

Definition 9 addresses both of the issues we mentioned above by, on one hand giving the distinguisher oracle access to v , and on the other hand by requiring that v computes an equivalence relation. The requirement that for any key-pair \mathbf{pk}, \mathbf{sk} and any pair of ciphertexts c, c' , if $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$ then $\delta(D_{\mathbf{sk}}(c), D_{\mathbf{sk}}(c')) \leq \text{negl}(k)$ is meant to capture the fact that the two ciphertexts c and c' can only be considered as replays of one another if they “carry essentially the same information”.

Definition 9 is written for a PKE scheme $\Pi = (G, E, D)$, but by taking the algorithm v required to exist by Definition 9 as a replay-filtering algorithm, we get a PKESF scheme $\Pi' = (G, E, D, v)$. Conversely, a PKESF scheme $\Pi = (G, E, D, F)$ is IND-srp-RCCA secure if the underlying PKE scheme $\Pi' = (G, E, D)$ is IND-srp-RCCA secure with respect to the filtering algorithm F of Π . Correctness of an IND-srp-RCCA secure PKESF Π' then follows from the correctness of the corresponding PKE $\Pi = (G, E, D)$.

It is instructive to see why Definition 9 does indeed require the filtering algorithm v to be meaningful. Consider, e.g., a trivial filtering algorithm such as the one that always sets $v(\mathbf{sk}, \mathbf{sk}, c, c') = 0$. This algorithm will not satisfy the definition above. But more importantly, it turns out that the above definition implies that Benchmark 2 is satisfied (see Theorem 3 further below), and by definition, Benchmark 2 requires the filtering algorithm to be meaningful (as otherwise the real and ideal systems are trivially distinguishable).

Lemma 4. *Consider any correct PKE scheme $\Pi = (G, E, D)$ that is IND-srp-RCCA secure, and let v be an algorithm with respect to which Π is IND-srp-RCCA secure. Then, $\Pi' = (G, E, D, v)$ is a correct PKESF scheme.*

Proof. The event $D_{\mathbf{sk}}(E_{\mathbf{pk}}(m)) \neq m \vee v(\mathbf{pk}, \mathbf{sk}, E_{\mathbf{pk}}(m), E_{\mathbf{pk}}(m')) = 1$ can only occur if at least one of $D_{\mathbf{sk}}(E_{\mathbf{pk}}(m)) \neq m$ or $v(\mathbf{pk}, \mathbf{sk}, E_{\mathbf{pk}}(m), E_{\mathbf{pk}}(m')) = 1$ occurs. From the correctness of Π , it follows that $D_{\mathbf{sk}}(E_{\mathbf{pk}}(m)) \neq m$ only occurs with negligible probability. Thus, it now only remains to show that $v(\mathbf{pk}, \mathbf{sk}, E_{\mathbf{pk}}(m), E_{\mathbf{pk}}(m')) = 1$ occurs with negligible probability.

Letting $c = E_{\mathbf{pk}}(m)$ and $c' = E_{\mathbf{pk}}(m')$, from the correctness of Π we have that $\delta(m, D_{\mathbf{sk}}(c)) \leq \text{negl}(k)$ and $\delta(m', D_{\mathbf{sk}}(c')) \leq \text{negl}(k)$. From the definition of IND-srp-RCCA security we have that if $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$ then $\delta(D_{\mathbf{sk}}(c), D_{\mathbf{sk}}(c')) \leq$

$\text{negl}(k)$. Combining these last 3 inequalities with the triangle inequality we find that $\delta(m, m') \leq \text{negl}(k)$. But note that m and m' are deterministic values (unlike $D_{\text{sk}}(c)$ and $D_{\text{sk}}(c')$ which are random variables over the distribution of the encryption and decryption randomness), hence we must have $\delta(m, m') = 0$ and $m = m'$. Putting this together, we have just shown that if $v(\text{pk}, \text{sk}, E_{\text{pk}}(m), E_{\text{pk}}(m')) = 1$ then $m = m'$.

Now, suppose that for some $m \in \mathcal{M}$ we have that with non-negligible probability $v(\text{pk}, \text{sk}, E_{\text{pk}}(m), E_{\text{pk}}(m)) = 1$ (i.e. v declares two fresh encryptions of the m as related). Then it is easy to create an efficient distinguisher \mathbf{D} that has non-negligible advantage in distinguishing the two IND-srp-RCCA game systems of Π with respect to v : First, \mathbf{D} makes a challenge query (`test messages`, m, \bar{m}) to the game system (where $m \neq \bar{m}$), and then \mathbf{D} generates a fresh encryption $c = E_{\text{pk}}(m)$ of m , and asks for the decryption of c to the game system. If the game system replies `test`, then \mathbf{D} outputs 0, and otherwise outputs 1. It is easy to see that \mathbf{D} 's advantage in distinguishing the two game systems is at least half of the probability that event $v(\text{pk}, \text{sk}, E_{\text{pk}}(m), E_{\text{pk}}(m)) = 1$ occurs, which by our assumption is non-negligible. Thus, \mathbf{D} has non-negligible advantage in distinguishing the two game systems, contradicting that Π is IND-srp-RCCA secure with respect to v . From this contradiction, it follows that for any m , $v(\text{pk}, \text{sk}, E_{\text{pk}}(m), E_{\text{pk}}(m)) = 1$ can only occur with negligible probability. \square

We now show that IND-srp-RCCA security of a PKESF $\Pi = (G, E, D, F)$ suffices for constructing an **RP-CONF** $[n]$ channel, i.e., satisfying Benchmark 2. To prove this, we create a simulator σ which internally uses F such that any distinguisher \mathbf{D} for $\text{enc}^A \text{dec}^B[\text{AUT}[1], \text{INS}[n], \mathbf{M}]$ and $\sigma^E \text{RP-CONF}[n]$ can be transformed into an equally good distinguisher for the IND-srp-RCCA game systems.

Theorem 3. *Let $\Pi = (G, E, D, F)$ be a correct PKESF scheme that is IND-srp-RCCA secure. There exists an efficient simulator σ and for any $n \in \mathbb{N}$ there exists an efficient reduction \mathbf{R} such that for every distinguisher \mathbf{D} ,*

$$\begin{aligned} \Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\text{AUT}[1], \text{INS}[n], \mathbf{M}], \sigma^E \text{RP-CONF}[n]) \\ = n \cdot \Delta^{\mathbf{DR}}(\mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{G}_1^{\text{IND-srp-RCCA}}). \end{aligned}$$

Proof. Consider the following simulator σ (defined in the Appendix, Algorithm 2) for interface E of **RP-CONF** $[n]$, which has two sub-interfaces denoted by $E.1$ and $E.2$ on the outside: Initially, the simulator generates a key-pair (pk, sk) , initializes two tables (`generatedCtxts` and `storedCtxts`) to empty, and outputs pk at $E.1$. When it receives the i -th input l_i , σ generates an encryption $\tilde{c} \leftarrow E_{\text{pk}}(\tilde{m})$ of a randomly chosen message \tilde{m} of length l_i , records in the table `generatedCtxts` the pair (i, \tilde{c}) and then outputs \tilde{c} at $E.2$. When c' is input at $E.2$, σ looks up the `storedCtxts` table for any ciphertext c that has been previously processed such that $F(\text{pk}, \text{sk}, c, c') = 1$. If the simulator finds one, then it simply ignores c' . Otherwise, it stores c' into the `storedCtxts` table, and then looks up the `generatedCtxts` table for any pair (i, c'') such that $F(\text{pk}, \text{sk}, c'', c') = 1$. If it finds

any such pair, σ outputs (dlv, j) to the **RP-CONF** $[n]$ channel (where j is the smallest index such that (j, c') is such that $F(\text{pk}, \text{sk}, c', c') = 1$). Otherwise, it decrypts c' , obtaining some plaintext m' , and then outputs (inj, m') into the **RP-CONF** $[n]$ channel (unless $m' = \perp$).

A distinguisher **D** connected to $\text{enc}^A \text{dec}^B[\text{AUT}[1], \text{INS}[n], \mathbf{M}]$ initially sees a public key at interface $E.1$. If **D** inputs a message m at interface A , an encryption of m is output at interface $E.2$. When **D** inputs a ciphertext c' at E , a decryption m' of c' is output at interface B , unless a F -related ciphertext has already been processed. When **D** is connected to the ideal $\sigma^E \text{RP-CONF}[n]$, it initially also sees a public key at $E.1$. But when it inputs the i -th message m_i at A , it sees an encryption c_i^* of a randomly chosen message m_i^* (where $|m_i^*| = |m_i|$) output at interface $E.2$. When some c' is input at interface $E.2$, if c' is F -related to some c input earlier then it is ignored and nothing happens. Otherwise, if a ciphertext c_i^* output earlier by σ is F -related to c' (where i is the smallest index satisfying this condition), then m_i is output at the B interface, unless it has already been processed by the **RP-CONF** $[n]$ channel, in which case nothing happens. Else, m' (the decryption of c') is output at interface B (unless $m' = \perp$).

We now introduce n reduction systems $\mathbf{R}_1, \dots, \mathbf{R}_n$, such that

S.1 $\mathbf{R}_1 \mathbf{G}_0^{\text{IND-srp-RCCA}} \equiv \text{enc}^A \text{dec}^B[\text{AUT}[1], \text{INS}[n], \mathbf{M}]$;

S.2 $\mathbf{R}_n \mathbf{G}_1^{\text{IND-srp-RCCA}} \equiv \sigma^E \text{RP-CONF}[n]$; and

S.3 for each $i \in \{1, \dots, n-1\}$ we have $\mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}} \equiv \mathbf{R}_{i+1} \mathbf{G}_0^{\text{IND-srp-RCCA}}$.

For $i = 1, \dots, n$, consider the following reduction system \mathbf{R}_i (defined in the Appendix, Algorithm 3) which processes at most n inputs at the outside A -interface: First, \mathbf{R}_i starts the game system, forwarding the public key pk generated by the game to the $E.1$ -interface. In addition, it also initializes the variable counter to 0, and the tables generatedCtxts and storedCtxts to empty. When a message m is input at the A -interface of \mathbf{R}_i , the variable counter is incremented, and then \mathbf{R}_i proceeds as follows:

counter $< i$: \mathbf{R}_i chooses a message \tilde{m} of length $|m|$ uniformly at random, and computes $c \leftarrow E_{\text{pk}}(\tilde{m})$; then, \mathbf{R}_i records (m, c) in the generatedCtxts table;
counter $= i$: \mathbf{R}_i chooses a message \tilde{m} of length $|m|$ uniformly at random, and makes the challenge query (**test messages**, m, \tilde{m}) to the game, which replies with some ciphertext c ; then, \mathbf{R}_i records (m, c) in the generatedCtxts table;
counter $> i$: \mathbf{R}_i computes $c \leftarrow E_{\text{pk}}(m)$, and records the pair (m, c) in the generatedCtxts table.

In all cases, \mathbf{R}_i then outputs c at the outside $E.2$ -interface. When c' is input at the $E.2$ -interface, \mathbf{R}_i behaves as follows. If a F -related ciphertext c has been recorded in the storedCtxts table, \mathbf{R}_i ignores c' , and so nothing happens. Otherwise, \mathbf{R}_i records c' into the storedCtxts table. Next, it checks for the existence of a pair (m, c) in the generatedCtxts table such that $F(\text{pk}, \text{sk}, c, c') = 1$. If there any such pair, then m is output at the B interface. Otherwise, \mathbf{R}_i queries the game for the decryption of c' , which returns some m . If $m \neq \perp$ then \mathbf{R}_i outputs m at the B interface, and otherwise it does not take any action. With this, it is easy to see that indeed conditions **S.1**, **S.2** and **S.3** hold. To conclude the proof, consider the reduction \mathbf{R} , which first chooses an index $i \in \{1, \dots, n\}$ uniformly at random,

and then simply behaves as the reduction system \mathbf{R}_i . Using the above facts and a simple triangle equality it follows

$$\begin{aligned}
& \Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[n], \mathbf{M}], \sigma^E \mathbf{RP-CONF}[n]) \\
&= \Delta^{\mathbf{D}}(\mathbf{R}_1 \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_n \mathbf{G}_1^{\text{IND-srp-RCCA}}) \\
&= \sum_{i=1}^n \left(\Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}) \right. \\
&\quad \left. + \Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}, \mathbf{R}_{i+1} \mathbf{G}_0^{\text{IND-srp-RCCA}}) \right) \\
&= \sum_{i=1}^n \Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}) \\
&= n \cdot \Delta^{\mathbf{DR}}(\mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{G}_1^{\text{IND-srp-RCCA}}). \quad \square
\end{aligned}$$

5.3 Achieving Benchmark 3: Constructing the RP-CONF Channel with Outsourceable Replay Protection

We now look into the setting where a third party who does not possess the secret-key is responsible for filtering out ciphertext replays (Benchmark 3). In this setting IND-srp-RCCA security seems too weak, as the algorithm v which the IND-srp-RCCA game systems use for detecting ciphertext replays (i.e. to check if two ciphertexts are replays of one another) have access to the secret-key. For this reason, we will now introduce the IND-prp-RCCA security notion, which is the analogous of IND-srp-RCCA security for public detection of ciphertext replays.

Definition 10. *A scheme $\Pi = (G, E, D)$ is IND-prp-RCCA secure if there is an efficient algorithm v that computes, for each public key \mathbf{pk} , an equivalence relation over ciphertexts c, c' such that for every \mathbf{pk} in the support of $G(1^k)$ and every pair of ciphertexts c, c' , if $v(\mathbf{pk}, c, c') = 1$ then $\delta(D_{\mathbf{sk}}(c), D_{\mathbf{sk}}(c')) \leq \text{negl}(k)$ (where the randomness is over the internal randomness of D and over the conditional distribution of the secret key \mathbf{sk} for the given public key \mathbf{pk} according to the key-pair distribution of $G(1^k)$), and if no efficient distinguisher \mathbf{D} distinguishes the two IND-prp-RCCA game systems (described ahead) with non-negligible advantage (in the security parameter k) over random guessing. The IND-prp-RCCA game systems work just as the IND-srp-RCCA game systems, except that now the game system does not have to provide the distinguisher with oracle access to v , as the distinguisher can anyway check whether any two ciphertexts are related according to v by itself.*

Even though IND-pd-RCCA security was introduced to capture efficient public detectability of ciphertext replays [7], and apart from the technical issues we already identified with its definition, IND-pd-RCCA still seems too weak to actually capture this feature for the replay detection algorithm is not required to compute an equivalence relation over ciphertexts.

Just like for IND-srp-RCCA, Definition 10 is written for a PKE scheme $\Pi = (G, E, D)$, but by taking the algorithm v required to exist by Definition 10 as a

replay-filtering algorithm, we get a PKEPF scheme $\Pi' = (G, E, D, v)$. Correctness of a IND-prp-RCCA secure PKEPF Π' then follows from the correctness of the corresponding PKE $\Pi = (G, E, D)$.

Lemma 5. *Consider any correct PKE scheme $\Pi = (G, E, D)$ that is IND-prp-RCCA secure, and let v be an algorithm with respect to which Π is IND-prp-RCCA secure. Then, $\Pi' = (G, E, D, v)$ is a correct PKEPF scheme.*

We omit the proof of Lemma 5 as it resembles the one of Lemma 4.

We now show that IND-prp-RCCA security of a PKEPF $\Pi = (G, E, D, F)$ suffices for constructing an **RP-CONF** $[n]$ channel even when the filtering is run by a third-party without access to the secret key, i.e., it satisfies Benchmark 3. To prove this, we have to show how to create a simulator σ which internally uses F such that any distinguisher \mathbf{D} for $\text{enc}^A \text{dec}^B \text{rp}^C[\mathbf{AUT}[1]_B, \mathbf{AUT}_C, \mathbf{INS}[n], \mathbf{M}]$ and $\sigma^E \mathbf{RP-CONF}[n]$ can be transformed into an equally good distinguisher for the IND-prp-RCCA game systems.

Theorem 4. *Let $\Pi = (G, E, D, F)$ be a correct and IND-prp-RCCA secure PKEPF scheme. There exists an efficient simulator σ and for any $n \in \mathbb{N}$ there exists an efficient reduction \mathbf{R} such that for every distinguisher \mathbf{D} ,*

$$\begin{aligned} \Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B \text{rp}^C[\mathbf{AUT}[1]_B, \mathbf{AUT}_C, \mathbf{INS}[n], \mathbf{M}], \sigma^E \mathbf{RP-CONF}[n]) \\ = \Delta^{\mathbf{DR}}(\mathbf{G}_0^{\text{IND-prp-RCCA}}, \mathbf{G}_1^{\text{IND-prp-RCCA}}). \end{aligned}$$

It is easy to see that this result can be obtained along the lines of Theorem 3.

6 Relating the Security Games

In this section we prove all the implications and separations between the game-based security notions that are depicted in Fig. 1.

Lemma 6. IND-CCA-2 \Rightarrow IND-prp-RCCA.

Proof. Define v so that $v(\text{pk}, c, c') = 1$ if and only if $c = c'$. Note that v satisfies IND-prp-RCCA security, since if $v(\text{pk}, c, c') = 1$ then $\delta(D_{\text{sk}}(c), D_{\text{sk}}(c')) = 0$. \square

Lemma 7. IND-prp-RCCA \Rightarrow IND-srp-RCCA.

Proof. Any algorithm v that satisfies IND-prp-RCCA also satisfies IND-srp-RCCA security (where v ignores the secret key sk). \square

The proof of the following result can be found in the Appendix, Sect. B.

Lemma 8. *Any correct and IND-srp-RCCA secure PKE scheme Π is $[n]$ IND-cl-RCCA secure.*

Lemma 9. $[n]$ IND-cl-RCCA $\Rightarrow [n - 1]$ IND-cl-RCCA.

Proof. Any distinguisher for the $[n - 1]$ IND-cl-RCCA game systems is also a distinguisher for the $[n]$ IND-cl-RCCA systems with the same advantage. \square

Lemma 10. $[1]$ IND-cl-RCCA \Rightarrow IND-RCCA.

Proof. From any distinguisher \mathbf{D} for the IND-RCCA game systems we create a distinguisher \mathbf{D}' for the $[1]$ IND-cl-RCCA game systems: \mathbf{D}' uses \mathbf{D} internally forwarding every query between \mathbf{D} and the $[1]$ IND-cl-RCCA game, except for decryption queries, where it behaves as follows: If, after the challenge plaintexts m_0 and m_1 are set, \mathbf{D} makes a decryption query of some ciphertext such that the $[1]$ IND-cl-RCCA game replies with either m_0 or m_1 , then \mathbf{D}' sends `test` to \mathbf{D} , and otherwise it sends what was output by the IND-RCCA game system. \square

Lemma 11. IND-RCCA $\not\Rightarrow$ $[1]$ IND-cl-RCCA.

Proof. By Theorem 2, $[1]$ IND-cl-RCCA security suffices for achieving Benchmark 1 for a single message. By Theorem 5, IND-RCCA does not suffice for achieving Benchmark 1 for a single message. \square

For the sake of simplicity, the two following results (Lemmata 12 and 13) assume the existence of an IND-CCA-2 secure PKE scheme. We note that both results can be generalized to only assume an $[n]$ IND-cl-RCCA (IND-srp-RCCA, respectively) secure scheme at the price of having a less elegant proof.

Lemma 12. $[n]$ IND-cl-RCCA $\not\Rightarrow$ IND-srp-RCCA.

Proof. From a IND-CCA-2 secure scheme $\Pi = (G, E, D)$, we create a scheme $\Pi' = (G', E', D')$ that is $[n]$ IND-cl-RCCA secure but not IND-srp-RCCA secure. Π' works just as Π except that now during encryption E' appends a bit 0 to the ciphertexts generated by E , and during decryption, if the last bit of the ciphertext is 0 then D' ignores it and decrypts the ciphertext using D , and otherwise, with $\frac{1}{2}$ probability D' outputs \perp and with the remaining $\frac{1}{2}$ probability D' ignores the last bit and decrypts the ciphertext using D .

Clearly, it is easy to create an algorithm v that suits $[n]$ IND-cl-RCCA such that no distinguisher has non-negligible advantage in distinguishing the two $[n]$ IND-cl-RCCA game systems for Π' with respect to v : for $b \in \{0, 1\}$, $v(\mathbf{pk}, \mathbf{sk}, c \parallel 0, c' \parallel b) = 1$ if and only if $c = c'$. On the other hand, any algorithm v' that suits IND-srp-RCCA cannot relate ciphertexts $c \parallel 0$ and $c \parallel 1$ since $\delta(D'_{\mathbf{sk}}(c \parallel 0), D'_{\mathbf{sk}}(c \parallel 1))$ is not negligible anymore. As such, a distinguisher can ask for the decryption of $c \parallel 1$ and use this to distinguish the game systems. \square

Lemma 13. IND-srp-RCCA $\not\Rightarrow$ IND-prp-RCCA.

Proof. From a IND-CCA-2 secure scheme $\Pi = (G, E, D)$, we create a scheme $\Pi' = (G', E', D')$ that is IND-srp-RCCA secure but not IND-prp-RCCA secure. Π' works just as Π except that now G' additionally picks a bit b uniformly at random and sets the key-pair to be $(\mathbf{pk}, (\mathbf{sk}, b))$, where $(\mathbf{pk}, \mathbf{sk})$ was the key-pair generated by G . More, during encryption E' uses E internally to generate a

ciphertext c and outputs (c, c) as the ciphertext, and during decryption, on input (c_0, c_1) , D' uses D internally to decrypt c_b (where b is the bit of the secret key that was sampled by G').

It is easy to create an algorithm v that suits IND-srp-RCCA such that no distinguisher has non-negligible advantage in distinguishing the two IND-srp-RCCA game systems for Π' with respect to v : for $b \in \{0, 1\}$, $v(\mathbf{pk}, \mathbf{sk}, (c_0, c_1), (c_0', c_1')) = 1$ if and only if $c_b = c_b'$, where b is again the bit of the secret key.

On the other hand, any algorithm v' that suits IND-prp-RCCA cannot relate ciphertext (c, c) with any of the following ciphertexts: (c, c_0') , (c, c_1') , (c_0', c) and (c_1', c) , where c_0' and c_1' are fresh encryptions of 0 and 1 respectively. This is so since, otherwise, either one could use v' to break the semantic security of Π (contradicting that it is IND-CCA-2 secure), or v' would not be suitable for IND-prp-RCCA, as one of $\delta(D'_{\mathbf{sk}}(c, c), D'_{\mathbf{sk}}(c, c_0'))$, $\delta(D'_{\mathbf{sk}}(c, c), D'_{\mathbf{sk}}(c, c_1'))$, $\delta(D'_{\mathbf{sk}}(c_0', c), D'_{\mathbf{sk}}(c_0', c))$ and $\delta(D'_{\mathbf{sk}}(c_1', c), D'_{\mathbf{sk}}(c_1', c))$ is not negligible anymore. As such, a distinguisher can ask for the decryption of these four ciphertexts and use the outputs to distinguish the IND-prp-RCCA game systems. \square

Lemma 14. IND-prp-RCCA $\not\Rightarrow$ IND-CCA-2.

Proof. Consider an IND-prp-RCCA secure PKE scheme $\Pi = (G, E, D)$; we create a scheme $\Pi' = (G', E', D')$ that is IND-prp-RCCA secure but not IND-CCA-2 secure: Π' works exactly as Π except that E' appends a bit 0 to the ciphertexts generated by E , and during decryption D' ignores the last bit added by E' is ignored. Since Π is IND-prp-RCCA secure, so is Π' . However, Π' is not IND-CCA-2 secure. \square

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A IND-RCCA Is Not Sufficient for Benchmark 1

In this section we prove that IND-RCCA security is, in general, too weak for achieving Benchmark 1 (see Sect. 3.1). As already mentioned, this result is in spirit with the separation proven in [7] between UC-RCCA and IND-RCCA for small message spaces.

Theorem 5. *There is an IND-RCCA secure PKE scheme Π for which there is an efficient distinguisher \mathbf{D} such that for any simulator σ ,*

$$\Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]], \sigma^E \mathbf{CONF}[1]) \geq \frac{1}{2}.$$

At a high level, we construct an IND-RCCA secure PKE scheme Π' for the binary message space that is malleable, in that an adversary can tamper a ciphertext into another that decrypts to a related message. While such tampering attacks do not help an adversary winning the IND-RCCA game for Π' ¹², we show that Benchmark 1 cannot be achieved using Π' , as it still allows an attacker to tamper with what Alice sends.

Let $\Pi = (G, E, D)$ be a correct and IND-RCCA secure PKE scheme for the binary message $\mathcal{M} = \{0, 1\}$. From Π , we construct a PKE scheme $\Pi' = (G', E', D')$, which works just as Π , except that now, E' appends an extra bit 0 to the ciphertexts, and during decryption D' uses D internally to decrypt the input ciphertext (ignoring the last bit appended by E'), and then XORs the plaintext output by D with the extra bit that was appended to the ciphertext during encryption (unless D outputs \perp , in which case D' also outputs \perp). It is easy to see that Π' is also IND-RCCA secure:

¹² Note that, even if the adversary manages to tamper the challenge ciphertext into one that decrypts to a different plaintext, it cannot leverage this attack into distinguishing the two game systems, because in the case of the binary message space the IND-RCCA game systems will not decrypt a ciphertext that decrypts to any of the two challenge plaintexts.

Lemma 15. *If Π is correct and IND-RCCA secure, then so is Π' .*

Proof. First, note that the correctness of Π' follows trivially from the correctness of Π . Now, consider a distinguisher \mathbf{D} for the IND-RCCA game systems of Π' . We create a distinguisher \mathbf{D}' for the IND-RCCA game systems of Π that distinguishes the two systems with the same advantage: First, \mathbf{D}' forwards the public key output by the game system to \mathbf{D} . Whenever \mathbf{D} makes a decryption query $(\text{ciphertext}, c \parallel \text{mask})$, \mathbf{D}' queries the game system on $(\text{ciphertext}, c)$. If the game system replies with some plaintext m , \mathbf{D}' sends $m \oplus \text{mask}$ back to \mathbf{D} as the output of the decryption query. Otherwise, if the game system replies with \perp , then \mathbf{D}' simply forwards \perp back to \mathbf{D} . When \mathbf{D} issues the challenge query $(\text{test messages}, m_0, m_1)$, \mathbf{D}' forwards it to the game system; when \mathbf{D}' receives the challenge ciphertext c^* back from the game system, it sends $c^* \parallel 0$ as the challenge ciphertext to \mathbf{D} . After the challenge ciphertext is set, whenever \mathbf{D} issues a decryption query $(\text{ciphertext}, c \parallel \text{mask})$, \mathbf{D}' queries the game system on $(\text{ciphertext}, c)$. If $m_0 \neq m_1$, or, if $m_0 = m_1$ and $\text{mask} = 0$, then \mathbf{D}' simply forwards the game system's reply back to \mathbf{D} . Otherwise, if $m_0 = m_1$ and $\text{mask} = 1$, \mathbf{D}' replies to \mathbf{D} :

- `test` if the game system output a plaintext;
- $m_0 \oplus 1$ if the game system output `test`; and
- \perp if the game system replied \perp .

Finally, \mathbf{D}' outputs as its guess whatever \mathbf{D} outputs as its guess. It is easy to see that \mathbf{D}' provides \mathbf{D} with a perfect emulation of the IND-RCCA game systems for Π' . It follows

$$\Delta^{\mathbf{D}'}(\mathbf{G}_0^{\Pi\text{-IND-RCCA}}, \mathbf{G}_1^{\Pi\text{-IND-RCCA}}) = \Delta^{\mathbf{D}}(\mathbf{G}_0^{\Pi'\text{-IND-RCCA}}, \mathbf{G}_1^{\Pi'\text{-IND-RCCA}}).$$

□

With this, we finally show that if Π' is used as the underlying PKE scheme for $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]]$, then Benchmark 1 is not achieved. More concretely, we show that there is a distinguisher \mathbf{D} such that for any simulator σ , \mathbf{D} distinguishes $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]]$ and $\sigma^E \mathbf{CONF}[1]$, where protocol $\pi = (\text{enc}, \text{dec})$ uses Π' as the underlying PKE scheme.

Lemma 16. *There is a distinguisher \mathbf{D} such that, for any possible simulator σ , $\Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]], \sigma^E \mathbf{CONF}[1]) \geq \frac{1}{2}$.*

Proof. Distinguisher \mathbf{D} behaves as follows: first, it chooses a message \tilde{m} uniformly at random from the message space $\mathcal{M} = \{0, 1\}$. Then, it inputs \tilde{m} into the channel's interface A . If afterwards either no ciphertext is output at interface E or the ciphertext which is output ends with bit 1, then \mathbf{D} can immediately output 1 as its guess (meaning that it is interacting with $\sigma^E \mathbf{CONF}[1]$), since the real world system would never behave like this. Otherwise, upon receiving a ciphertext $c \parallel 0$, \mathbf{D} inputs the ciphertext $c \parallel 1$ into the same interface. Finally, if $\tilde{m} \oplus 1$ is output at the B interface, then \mathbf{D} guesses that it is interacting with

$\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]]$ outputting 0 as its guess. Otherwise, it guesses that it is interacting with $\sigma^E \mathbf{CONF}[1]$ and outputs 1.

First, note that if no ciphertext is output at interface E or the ciphertext which is output ends with bit 1, then \mathbf{D} guesses correctly with probability 1. As such from now on, we assume that a ciphertext $c \parallel 0$ is output at interface E after \mathbf{D} inputs \tilde{m} into interface A . Note that, when \mathbf{D} interacts with $\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]]$, the PKE scheme Π' allows \mathbf{D} to tamper the ciphertext into one that decrypts to a different, but related, message. However, due to the definition of the ideal \mathbf{CONF} channel, it is impossible for \mathbf{D} to mount an analogous attack — where \mathbf{D} manages to tamper with whatever Alice sends — when \mathbf{D} is interacting with $\sigma^E \mathbf{CONF}[1]$. Consequently, regardless of whichever simulator σ one connects to the ideal \mathbf{CONF} channel, it is impossible for σ to translate such attack from the real world construction into the ideal \mathbf{CONF} channel: When \mathbf{D} injects the ciphertext $c \parallel 1$ into σ , σ can either choose to issue a $(\text{dlv}, 1)$ query to the ideal \mathbf{CONF} channel (forwarding whatever was input at the A interface to the B interface), or it can choose to issue a (inj, m) query (making the ideal \mathbf{CONF} channel output m at the B interface), for some plaintext m . However, if σ issues a $(\text{dlv}, 1)$ query, \mathbf{D} immediately notices that it is interacting with the ideal \mathbf{CONF} channel, as its attack had no effect. On the other hand, if σ chooses to issue a (inj, m) query, then with probability $\frac{1}{2}$, the injected m will be such that $m = \tilde{m}$. Consequently, with probability $\frac{1}{2}$, \mathbf{D} still learns that it is interacting with the ideal world (and thus outputs 1). Hence, for any simulator σ , $\Delta^{\mathbf{D}}(\text{enc}^A \text{dec}^B[\mathbf{AUT}[1], \mathbf{INS}[1]], \sigma^E \mathbf{CONF}[1]) \geq \frac{1}{2} - 0 = \frac{1}{2}$. \square

From Lemmata 15 and 16 it follows that IND-RCCA security is not sufficient for achieving Benchmark 1, thus concluding the proof of Theorem 5.

B Proof of Lemma 8

Proof (Lemma 8). To prove the claim, we first introduce the notion of $[n]$ IND-srp-RCCA security: A scheme $\Pi = (G, E, D)$ is $[n]$ IND-srp-RCCA secure if it is secure according to the definition of IND-srp-RCCA security, except that now the $[n]$ IND-srp-RCCA game systems work as follows (for $b \in \{0, 1\}$):

Initialization: First, $\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ creates a table t of ciphertexts which is initially empty. Then, it runs $(\mathbf{pk}, \mathbf{sk}) \leftarrow G(1^k)$, and sends \mathbf{pk} to \mathbf{D} .

Decryption queries: Whenever \mathbf{D} queries $(\text{ciphertext}, c)$, the game system calls, for each ciphertext c_j^* stored in t , $v(\mathbf{pk}, \mathbf{sk}, c_j^*, c)$. If for every ciphertext in t v 's output is 0, then the game system replies with the decryption $m = D_{\mathbf{sk}}(c)$ of c to \mathbf{D} . Otherwise, let c_k^* be the ciphertext stored in t with the smallest k such that $v(\mathbf{pk}, \mathbf{sk}, c_k^*, c) = 1$. Then, $\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ replies (test, k) to \mathbf{D} .

i -th challenge query (for $i \leq n$): Whenever the distinguisher \mathbf{D} issues a challenge query $(\text{test messages}, m_{0,i}, m_{1,i})$, where $m_{0,i}, m_{1,i} \in \mathcal{M}$ such that $|m_{0,i}| = |m_{1,i}|$, the game system computes $c_i^* = E_{\mathbf{pk}}(m_{b,i})$, stores c_i^* in table t , and sends c_i^* to \mathbf{D} .

As for IND-srp-RCCA security, the $[n]$ IND-srp-RCCA game systems provide \mathbf{D} with oracle access to v throughout the entire game, so \mathbf{D} can check whether any two ciphertexts c, c' are related according to v with respect to the key-pair \mathbf{pk}, \mathbf{sk} generated by the game system.

Having introduced $[n]$ IND-srp-RCCA security, we now prove that IND-srp-RCCA security implies $[n]$ IND-srp-RCCA security. To conclude the proof, we will show $[n]$ IND-srp-RCCA $\Rightarrow [n]$ IND-cl-RCCA.

Consider a scheme Π which is IND-srp-RCCA secure with respect to some algorithm v . Thus, no efficient distinguisher has non-negligible advantage in distinguishing the two IND-srp-RCCA game systems for Π with respect to v . Note that, since v suits IND-srp-RCCA, it also suits $[n]$ IND-srp-RCCA. We now construct, from any distinguisher \mathbf{D} for distinguishing the $[n]$ IND-srp-RCCA game systems, a distinguisher \mathbf{DR} (where \mathbf{R} is an efficient reduction) for distinguishing the IND-srp-RCCA game systems, such that

$$\Delta^{\mathbf{DR}}(\mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{G}_1^{\text{IND-srp-RCCA}}) = \frac{1}{n} \cdot \Delta^{\mathbf{D}}(\mathbf{G}_0^{[n]\text{IND-srp-RCCA}}, \mathbf{G}_1^{[n]\text{IND-srp-RCCA}}).$$

To this end, we introduce n reduction systems $\mathbf{R}_1, \dots, \mathbf{R}_n$, such that

- S.1** $\mathbf{R}_1 \mathbf{G}_0^{\text{IND-srp-RCCA}} \equiv \mathbf{G}_0^{[n]\text{IND-srp-RCCA}}$,
- S.2** $\mathbf{R}_n \mathbf{G}_1^{\text{IND-srp-RCCA}} \equiv \mathbf{G}_1^{[n]\text{IND-srp-RCCA}}$; and
- S.3** for each $i \in \{1, \dots, n-1\}$ we have $\mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}} \equiv \mathbf{R}_{i+1} \mathbf{G}_0^{\text{IND-srp-RCCA}}$.

For $i = 1, \dots, n$, consider the following reduction system \mathbf{R}_i (which processes at most n challenge queries): Initially, \mathbf{R}_i forwards the public key \mathbf{pk} generated by the game system to \mathbf{D} . When \mathbf{D} submits the j -th challenge query (**test messages**, $m_{0,j}, m_{1,j}$) to \mathbf{R}_i : 1. if $j < i$, \mathbf{R}_i computes $c_j^* \leftarrow E_{\mathbf{pk}}(m_{1,j})$ and then records the pair (j, c_j^*) ; 2. if $j = i$, \mathbf{R}_i makes the challenge query (**test messages**, $m_{0,j}, m_{1,j}$) to the game system, which replies with some ciphertext c_j^* and then it records the pair (j, c_j^*) ; finally, 3. if $j > i$, \mathbf{R}_i computes $c_j^* \leftarrow E_{\mathbf{pk}}(m_{0,j})$, and records the pair (j, c_j^*) . In all cases, \mathbf{R}_i then outputs c_j^* as the challenge ciphertext to \mathbf{D} . When \mathbf{D} makes a decryption query (**ciphertext**, c), and letting l denote the number of challenge queries already issued by \mathbf{D} , \mathbf{R}_i behaves as follows: \mathbf{R}_i iterates each pair (j, c^*) recorded (where $j = 1, \dots, l$) in ascending order, and when it finds a pair such that $v(\mathbf{pk}, \mathbf{sk}, c^*, c) = 1$, \mathbf{R}_i outputs (**test**, j). If it does not find any such pair, \mathbf{R}_i forwards the decryption query (**ciphertext**, c) to the game system, and forwards the game system's reply back to \mathbf{D} . It is easy to see that all three conditions (**S.1**, **S.2** and **S.3**) hold. To conclude the proof that IND-srp-RCCA security implies $[n]$ IND-srp-RCCA security, consider the reduction system \mathbf{R} , which first chooses an index $i \in \{1, \dots, n\}$ uniformly at random, and then simply behaves as the reduction system \mathbf{R}_i . Using

the above facts and a simple triangle equality it follows

$$\begin{aligned}
& \Delta^{\mathbf{D}}(\mathbf{G}_0^{[n]\text{IND-srp-RCCA}}, \mathbf{G}_1^{[n]\text{IND-srp-RCCA}}) \\
&= \Delta^{\mathbf{D}}(\mathbf{R}_1 \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_n \mathbf{G}_1^{\text{IND-srp-RCCA}}) \\
&= \sum_{i=1}^n \left(\Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}) \right. \\
&\quad \left. + \Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}, \mathbf{R}_{i+1} \mathbf{G}_0^{\text{IND-srp-RCCA}}) \right) \\
&= \sum_{i=1}^n \Delta^{\mathbf{D}}(\mathbf{R}_i \mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{R}_i \mathbf{G}_1^{\text{IND-srp-RCCA}}) \\
&= \frac{1}{n} \cdot \Delta^{\mathbf{DR}}(\mathbf{G}_0^{\text{IND-srp-RCCA}}, \mathbf{G}_1^{\text{IND-srp-RCCA}}).
\end{aligned}$$

To conclude, we now prove $[n]\text{IND-srp-RCCA} \Rightarrow [n]\text{IND-cl-RCCA}$. First, note that any algorithm v that suits $[n]\text{IND-srp-RCCA}$ security for some scheme Π also suits $[n]\text{IND-cl-RCCA}$ security for Π . Now, we show that any \mathbf{D} distinguishing the $[n]\text{IND-cl-RCCA}$ game systems for Π with respect to v can be transformed via an efficient reduction \mathbf{R} into a distinguisher \mathbf{DR} distinguishing the $[n]\text{IND-srp-RCCA}$ game systems for Π (also with respect to v) with essentially the same advantage. For $b \in \{0, 1\}$, \mathbf{DR} uses \mathbf{D} internally, forwarding everything back and forth between \mathbf{D} and the game system $\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$. Note that, \mathbf{R} , when connected to the game system $\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$, provides \mathbf{D} with a view that only negligibly differs from $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$: $\mathbf{R}\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ behaves just as $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ for every query that \mathbf{D} makes, except for decryption queries. When \mathbf{D} makes a decryption query ($\text{ciphertext}, c$), if $\mathbf{R}\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ replies with a decryption m of c , then $v(\text{pk}, \text{sk}, c^*, c) = 0$, implying that in this case $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ would also reply with m . Thus, in this case the two systems $\mathbf{R}\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ and $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ are perfectly indistinguishable. Now consider the case where $\mathbf{R}\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ replies (test, i). Then, $v(\text{pk}, \text{sk}, c^*, c) = 1$, implying $\delta(D_{\text{sk}}(c^*), D_{\text{sk}}(c)) \leq \text{negl}(k)$ (where k is the security parameter). From the correctness of Π (Definition 2), it follows that with high probability (i.e. with probability at least $1 - \text{negl}(k)$), the key-pair (pk, sk) is such that for every $m \in \mathcal{M}$, $\Pr[D_{\text{sk}}(E_{\text{pk}}(m)) \neq m] \leq \text{negl}(k)$. Thus, with high probability $\delta(D_{\text{sk}}(c^*), m_b) \leq \text{negl}(k)$ ¹³, implying $\delta(D_{\text{sk}}(c), m_b) \leq \text{negl}(k)$. Thus, $\mathbf{R}\mathbf{G}_b^{[n]\text{IND-srp-RCCA}}$ and $\mathbf{G}_b^{[n]\text{IND-cl-RCCA}}$ are perfectly indistinguishable with high probability, up to when $D_{\text{sk}}(c) \neq m_b$, which occurs with negligible probability. As such, \mathbf{DR} has essentially the same advantage in distinguishing $\mathbf{G}_0^{[n]\text{IND-srp-RCCA}}$ and $\mathbf{G}_1^{[n]\text{IND-srp-RCCA}}$ as \mathbf{D} has in distinguishing $\mathbf{G}_0^{[n]\text{IND-cl-RCCA}}$ and $\mathbf{G}_1^{[n]\text{IND-cl-RCCA}}$, concluding the proof. \square

¹³ Here, m_b denotes the probability distribution over $\mathcal{M} \cup \{\perp\}$ such that only m_b is in the support, thus occurring with probability 1.

C Pseudocode for Simulator and Reduction Systems of Theorem 3

In this section we present the pseudocode of the simulator σ (Algorithm 2) and reduction system \mathbf{R}_i (Algorithm 3) used for the proof of Theorem 3.

Algorithm 2 Simulator σ .

```

1: procedure INITIALIZATION
2:    $(\mathbf{pk}, \mathbf{sk}) \leftarrow G(1^k)$ 
3:   counter  $\leftarrow 0$ 
4:   generatedCtxts-Initialize( $\emptyset$ )
5:   storedCtxts-Initialize( $\emptyset$ )
6:    $E.1$ -Output( $\mathbf{pk}$ )
7: end procedure

8: procedure RP-CONF-INPUT( $l$ )
9:   counter  $\leftarrow$  counter + 1
10:   $\tilde{m} \leftarrow^{\mathcal{S}} \mathcal{M}^l$ 
11:   $\tilde{c} \leftarrow E_{\mathbf{pk}}(\tilde{m})$ 
12:  generatedCtxts-Add(counter,  $\tilde{c}$ )
13:   $E.2$ -Output( $\tilde{c}$ )
14: end procedure

15: procedure  $E.2$ -INPUT( $c'$ )
16:  if storedCtxts-Contains( $c$ ) s.t.  $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$  then
17:    Ignore  $c'$ 
18:  else
19:    storedCtxts-Add( $c'$ )
20:    if generatedCtxts-Contains( $i, c$ ) s.t.  $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$  then
21:      RP-CONF-Output(dlv,  $i$ )
22:    else
23:       $m' \leftarrow D_{\mathbf{sk}}(c')$ 
24:      if  $m' \neq \perp$  then
25:        RP-CONF-Output(inj,  $m'$ )
26:      end if
27:    end if
28:  end if
29: end procedure

```

Algorithm 3 Reduction \mathbf{R}_i .

```
1: procedure INITIALIZATION
2:    $\mathbf{pk} \leftarrow G.\text{Initialize}(1^k)$ 
3:   counter  $\leftarrow 0$ 
4:   generatedCtxts-Initialize( $\emptyset$ )
5:   storedCtxts-Initialize( $\emptyset$ )
6:    $E.1\text{-Output}(\mathbf{pk})$ 
7: end procedure

8: procedure A-INPUT( $m$ )
9:   counter  $\leftarrow$  counter + 1
10:  if counter  $< i$  then
11:     $\tilde{m} \leftarrow^{\$} \mathcal{M}^{|m|}$ 
12:     $c \leftarrow E_{\mathbf{pk}}(\tilde{m})$ 
13:    generatedCtxts-Add( $m, c$ )
14:  else if counter =  $i$  then
15:     $\tilde{m} \leftarrow^{\$} \mathcal{M}^{|m|}$ 
16:     $c \leftarrow G.\text{query}(\text{text messages}, m, \tilde{m})$ 
17:    generatedCtxts-Add( $m, c$ )
18:  else
19:     $c \leftarrow E_{\mathbf{pk}}(m)$ 
20:    generatedCtxts-Add( $m, c$ )
21:  end if
22:   $E.2\text{-Output}(\tilde{c})$ 
23: end procedure

24: procedure E.2-INPUT( $c'$ )
25:  if storedCtxts-Contains( $c$ ) s.t.  $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$  then
26:    Ignore  $c'$ 
27:  else
28:    storedCtxts-Add( $c'$ )
29:    if generatedCtxts-Contains( $m, c$ ) s.t.  $v(\mathbf{pk}, \mathbf{sk}, c, c') = 1$  then
30:       $B\text{-Output}(m)$ 
31:    else
32:       $m' \leftarrow G.\text{query}(\text{ciphertext}, c')$ 
33:      if  $m' \neq \perp$  then
34:         $B\text{-Output}(m')$ 
35:      else
36:        Ignore  $c'$ 
37:      end if
38:    end if
39:  end if
40: end procedure
```
