On the bijectivity of the map χ

Anna-Maurin Graner ^{*1}, Björn Kriepke ^{†1}, Lucas Krompholz ^{‡1}, and Gohar M. Kyureghyan^{§1}

¹Institute of Mathematics, University of Rostock, Germany

February 7, 2024

We prove that for n > 1 the map $\chi : \mathbb{F}_q^n \to \mathbb{F}_q^n$, defined by $y = \chi(x)$ with $y_i = x_i + x_{i+2} \cdot (1 + x_{i+1})$ for $1 \le i \le n$, is bijective if and only if q = 2 and n is odd, as it was conjectured in [8].

1 Introduction

Let q be any prime power and n a positive integer. Several cryptographic primitives, including ASCON [4] and SHA-3 [6], use the map $\chi : \mathbb{F}_q^n \to \mathbb{F}_q^n$ given by $y = \chi(x)$ with

$$y_i = x_i + x_{i+2} \cdot (1 + x_{i+1})$$

for $1 \leq i \leq n$, where the indices are computed modulo n. Let the symbol \odot denote the element wise multiplication of two vectors (also known as the Hadamard product), i.e., $z = x \odot y$ with $z_i = x_i \cdot y_i$ for all $i = 1, \ldots, n$. Further, denote by S the cyclic left shift operator on \mathbb{F}_q^n , that is $S(x_1, \ldots, x_n) = (x_2, \ldots, x_n, x_1)$. Let S^j denote the *j*-th iterate of S for $j \geq 0$. Note that S^0 is the identity map. Then χ can also be written as

$$\chi(x) = x + \mathcal{S}(x) \odot \mathcal{S}^2(x) + \mathcal{S}^2(x).$$

It is known that $\chi : \mathbb{F}_2^n \to \mathbb{F}_2^n$ is bijective if and only if *n* is odd [2]. Some partial results are proved about bijectivity of χ for $q \neq 2$. In [8] it was shown that for $k \geq 1$ the map χ is not a permutation, when

• q is odd,

^{*}anna-maurin.graner@uni-rostock.de

 $^{^{\}dagger} bjoern. kriepke@uni-rostock. de$

[‡]lucas.krompholz@uni-rostock.de

[§]gohar.kyureghyan@uni-rostock.de

- $q = 2^k$ and n is even,
- $q = 2^{2k}$ and n > 1 is odd,
- $q = 2^{3k}$ and n > 1 is odd.

In [7] the following additional parameters were ruled out using an approach based on Gröbner basis:

• $q = 2^{5k}$ or $q = 2^{7k}$ and n is a multiple of 3 or 5.

It was conjectured in [8] that χ is not a permutation in all other cases except when q = 2 and n odd. We confirm this conjecture using linear algebra methods. More precisely, we prove in Lemmas 3 to 5 that the following result holds:

Theorem 1. For q = 2 the map χ is a permutation if and only if n is odd. For any prime power q > 2, the map $\chi : \mathbb{F}_q^n \to \mathbb{F}_q^n$ is a permutation if and only if q is even and n = 1.

We conclude our note with a short proof for the rank of the linear part of $\chi(x+a)+\chi(x)$, which appears in the study of the differential properties of the map $\chi: \mathbb{F}_2^n \to \mathbb{F}_2^n$.

2 Deriving the linear system

The map χ is not a permutation if and only if there exist vectors $a, x \in \mathbb{F}_q^n$ with $a \neq 0$ such that

$$\chi(x+a) - \chi(x) = 0. \tag{1}$$

Note that for any j the map S^j is linear over \mathbb{F}_q . Furthermore, the Hadamard product is commutative and distributive with respect to addition, i.e. $x \odot y = y \odot x$ and $x \odot (y+z) = x \odot y + x \odot z$ for all $x, y, z \in \mathbb{F}_q^n$. Moreover, we have $S^j(x \odot y) = S^j(x) \odot S^j(y)$. Using these properties, we obtain

$$\begin{split} \chi(x+a) &= x + a + \mathbf{S}(x+a) \odot \mathbf{S}^2(x+a) + \mathbf{S}^2(x+a) \\ &= x + a + [\mathbf{S}(x) + \mathbf{S}(a)] \odot [\mathbf{S}^2(x) + \mathbf{S}^2(a)] + \mathbf{S}^2(x) + \mathbf{S}^2(a) \\ &= x + a + \mathbf{S}(x) \odot \mathbf{S}^2(x) + \mathbf{S}(x) \odot \mathbf{S}^2(a) + \mathbf{S}(a) \odot \mathbf{S}^2(x) + \mathbf{S}(a) \odot \mathbf{S}^2(a) + \mathbf{S}^2(a) \\ &= \chi(x) + a + \mathbf{S}(x) \odot \mathbf{S}^2(a) + \mathbf{S}(a) \odot \mathbf{S}^2(x) + \mathbf{S}(a) \odot \mathbf{S}^2(a) + \mathbf{S}^2(a) \\ \end{split}$$

and therefore

$$\chi(x+a) - \chi(x) = a + S^2(a) + S(a \odot S(x) + x \odot S(a) + a \odot S(a))$$

For a fixed $a \in \mathbb{F}_a^n \setminus \{0\}$, the equation $\chi(x+a) - \chi(x) = 0$ has a solution x if and only if

$$-a - S^{2}(a) = S(a \odot S(x) + x \odot S(a) + a \odot S(a))$$

has a solution, which, by applying S^{-1} on both sides, is equivalent to

$$-S^{-1}(a) - S(a) - a \odot S(a) = a \odot S(x) + x \odot S(a).$$
⁽²⁾

The right-hand side of (2) is a linear map in x and hence it reduces to a system of linear equations over \mathbb{F}_q . We represent this system of equations using matrices:

$$\begin{pmatrix} a_2 & a_1 & & & \\ & a_3 & a_2 & & & \\ & & a_4 & a_3 & & \\ & & \ddots & \ddots & & \\ & & & a_{n-1} & a_{n-2} & \\ & & & & & a_n & a_{n-1} \\ a_n & & & & & & a_1 \end{pmatrix} \cdot x = - \begin{pmatrix} a_1a_2 + a_2 + a_n \\ a_2a_3 + a_3 + a_1 \\ a_3a_4 + a_4 + a_2 \\ \vdots \\ a_{n-2}a_{n-1} + a_{n-1} + a_{n-3} \\ a_{n-1}a_n + a_n + a_{n-2} \\ a_na_1 + a_1 + a_{n-1} \end{pmatrix}, \quad (3)$$

where $a = (a_1, \ldots, a_n)$. We denote the coefficient matrix in (3) by A(a) and the vector in its right-hand side by b(a). We abbreviate $A(a) \cdot x = b(a)$ often by (A(a)|b(a)).

Observe that the map $\chi : \mathbb{F}_q^n \to \mathbb{F}_q^n$ is bijective if and only if for any non-zero $a \in \mathbb{F}_q^n$ equation (3) has no solution. Our goal is now to check whether (3) has a solution x for some fixed non-zero a.

3 The case q > 2

In this section we show that for q > 2 the map χ is a permutation on \mathbb{F}_q^n if and only if q is even and n = 1. We consider separately the cases n = 1, 2, 3 and n > 3.

First let us assume that n = 1. In that case S(x) = x is the identity map and therefore $\chi(x) = x + S(x) \odot S^2(x) + S^2(x) = x + x^2 + x = x^2 + 2x = x(x+2)$, which is a permutation if and only if q is even.

Remark 2. Note that for n = 1 in odd characteristic $\chi(0) = \chi(-2)$. In general for any n it holds that $\chi(0, \ldots, 0) = \chi(-2, \ldots, -2)$ and therefore χ is never a permutation in odd characteristic, as noted in [8]. Therefore, from now on we could restrict ourselves to even characteristic. However, the rest of the proof presented here is valid independently of the characteristic of \mathbb{F}_q , with the minor exception in the case n = 3.

We continue with n = 2. In this case (3) has the form

$$\left(\begin{array}{ccc} a_2 & a_1 & -a_1a_2 - 2a_2 \\ a_2 & a_1 & -a_1a_2 - 2a_1 \end{array}\right)$$

This has a solution for example in the case a = (1, 1) which shows that χ is not a permutation.

Next, let n = 3. Now the system (3) looks like

$$\left(\begin{array}{ccc|c} a_2 & a_1 & -a_1a_2 - a_2 - a_3 \\ a_3 & a_2 & -a_2a_3 - a_3 - a_1 \\ a_3 & a_1 & -a_3a_1 - a_1 - a_2 \end{array}\right).$$

Note that the determinant of the coefficient matrix is $2a_1a_2a_3$. Therefore, if q is odd, we can choose a_1, a_2, a_3 all nonzero and the corresponding system always has a solution. In the case q even, assume that $a_2 \neq 0$. Using the Gaussian elimination, we obtain

$$\begin{pmatrix} a_2 & a_1 & & & a_1a_2 + a_2 + a_3 \\ a_3 & a_2 & & & a_2a_3 + a_3 + a_1 \\ & & 0 & a_1^2 + a_2^2 + a_3^2 + a_1a_2 + a_1a_3 + a_2a_3 + a_1a_2a_3 \end{pmatrix}.$$
 (4)

This system has a solution if there exist choices of $a_1, a_2, a_3 \in \mathbb{F}_q$ such that $a_2, a_3 \neq 0$ and

$$a_1^2 + (a_2 + a_3 + a_2 a_3)a_1 + (a_2 a_3 + a_2^2 + a_3^2) = 0,$$
(5)

which is a quadratic equation in a_1 . Having in mind, that in binary fields a quadratic equation $X^2 + uX + v = 0$ has always a solution if u = 0, we put $a_2 + a_3 + a_2a_3 = 0$ in (5). Equivalently, by adding 1 on both sides, $(a_2 + 1)(a_3 + 1) = 1$. As q > 2, we can choose an element $a_3 \in \mathbb{F}_q \setminus \{0, 1\}$ and then $a_2 = \frac{1}{a_3+1} + 1 = \frac{a_3}{a_3+1} \neq 0$. For these $a_2, a_3 \neq 0$ the quadratic equation (5) has a solution $a_1 \in \mathbb{F}_q$, implying the existence of $(a_1, a_2, a_3) \neq 0$ for which the linear system (4) has a solution x.

We have thus proved the following lemma.

Lemma 3. Let q > 2. If n = 1 then χ is a permutation if and only if q is even. If n = 2, 3 then χ is not a permutation.

Let now n > 3. Again, we show that for certain choices of the vector $a \in \mathbb{F}_q^n \setminus \{0\}$ the equation (3) admits a solution x. Let $a_n = 0$. Then the linear system (3) reduces to

Further, let all a_1, \ldots, a_{n-1} be non-zero and assume

$$\det \begin{pmatrix} a_{n-1} & a_{n-2} \\ a_1 & a_1 + a_{n-1} \end{pmatrix} = 0,$$

or equivalently, $a_{n-1}(a_1 + a_{n-1}) = a_1 a_{n-2}$. Under this assumption there is a solution $x \in \mathbb{F}_q^n$. Indeed we can choose x_{n-1} arbitrarily, for example $x_{n-1} = 1$, and then $x_n = -\frac{a_{n-2}}{a_{n-1}}$. The remaining components are obtained by simple back substitution, as the other diagonal entries are all nonzero.

Now it remains to see that there are non-zero $a_1, a_{n-1}, a_{n-2} \in \mathbb{F}_q$ such that the assumption $a_{n-1}(a_1+a_{n-1}) = a_1a_{n-2}$ is satisfied. Note that because n > 3 the components

 a_1, a_{n-1}, a_{n-2} do not coincide. Let $a_{n-1} = 1$ and choose $a_1 \in \mathbb{F}_q \setminus \{0, -1\}$ arbitrarily. Then $a_1 + 1 \neq 0$ and $a_{n-2} = \frac{a_1 + 1}{a_1} \neq 0$, fulfilling the requirements.

We have thus proved the following result.

Lemma 4. Let q > 2 and n > 3. Then $\chi : \mathbb{F}_q^n \to \mathbb{F}_q^n$ is not a permutation.

4 The special case q = 2

It is known that for q = 2 the map $\chi : \mathbb{F}_2^n \to \mathbb{F}_2^n$ is bijective if and only if n is odd. If n is even it is easy to see that χ is not a permutation. Indeed,

$$\chi(1,0,1,0,\ldots,1,0) = (0,\ldots,0) = \chi(0,\ldots,0),$$

as it has been noted in [2]. The fact that χ is a permutation for n odd was proved in [2] by using a seed-and-leap method to compute the preimage of a given element $y \in \mathbb{F}_2^n$. A more detailed proof of this approach can be found in [3]. Another method to compute the inverse of χ for n odd is given in Appendix D of [1], however without a proof. In [5] an explicit inverse formula of χ is given and proved.

To have a unified proof for Theorem 1, we present here a short proof for the statement that $\chi : \mathbb{F}_2^n \to \mathbb{F}_2^n$ is bijective if n is odd, applying the method developed in the previous sections.

Let now n be odd. If n = 1 then $\chi(x) = x^2 = x$ which is a permutation. So now assume $n \ge 3$. Let $a \in \mathbb{F}_2^n \setminus \{0\}$ be arbitrary. We aim to show that there is no solution x to $\chi(x) + \chi(x+a) = 0$. It can be easily seen that χ is shift-invariant, i.e. $S(\chi(x)) = \chi(S(x))$ for all $x \in \mathbb{F}_2^n$. Therefore, if $\chi(x) + \chi(x+a) = 0$ has a solution, then it follows that also

$$0 = S(0) = S(\chi(x) + \chi(x+a)) = \chi(S(x)) + \chi(S(x) + S(a))$$

and there also exists a solution S(x) for S(a).

In the following we show that (3) has no solution by considering three cases. First we assume that a has two consecutive entries which are zero. Next we will assume that a has a zero entry such that the entries before and after are both nonzero. And finally we will assume that a only has nonzero entries.

Suppose now (3) has a solution x for a non-zero a with $a_i = a_{i+1} = 0$ for some $1 \le i \le n$. Since χ is shift-invariant, by considering an appropriate shift of a, we may assume without loss of generality that $a_n = a_1 = 0$. The last row of (3) then looks as follows:

$$\left(\begin{array}{ccc} 0 & & & 0 \mid a_{n-1} \end{array}\right).$$

As the system has a solution x, it then follows that $a_{n-1} = 0$. However, then by considering the (n-1)-th row, it follows that also $a_{n-2} = 0$. By repeating this argument we obtain a = 0, a contradiction.

Next we assume that there exists an index $i \in \{1, ..., n\}$ such that $a_i = 0$ and $a_{i-1} = a_{i+1} = 1$. Again, by considering shifts of a, we may assume that i = n. From

the last two rows of (3) it then immediately follows that $a_{n-2} = x_n = 0$. If $a_{n-3} = 0$, then we are in the previous case. Otherwise, we can repeat this argument and obtain that $a_{n-2k} = 0$ for all integers k. However, using that n = 2m + 1 is odd, we then also obtain $a_{n-2m} = a_1 = 0$, a contradiction to the assumption that $a_1 \neq 0$.

Finally, we need to consider a = (1, ..., 1). In this case (3) reduces to

By adding every of the first n-1 rows to the last one, we obtain (using that n-1 is even) the row

 $\left(\begin{array}{ccc} 0 & & 0 & | 1 \end{array} \right)$

which means that the equation has no solution.

The above considerations imply the following result:

Lemma 5. The map $\chi : \mathbb{F}_2^n \to \mathbb{F}_2^n$ is a permutation if and only if n is odd.

5 Rank of the coefficient matrix A(a) over \mathbb{F}_2

The equation (3) appears in the study of differential and linear properties of χ . In particular, the ranks of matrices A(a) allow to determine the Walsh spectrum of χ . In [8] the following proposition is proved:

Proposition 6. For any $a \in \mathbb{F}_2^n$ the rank of the matrix A(a) over \mathbb{F}_2 is given by

$$\operatorname{rank} A(a) = \omega(a) := \begin{cases} n-1, & a = (1, \dots, 1) \\ \operatorname{wt}(a) + \operatorname{r}(a), & otherwise \end{cases}$$

where wt(a) is the Hamming weight and r(a) is the number of 001-patterns in a. More precisely, r(a) is the number of indices i = 1, ..., n such that $(a_i, a_{i+1}, a_{i+2}) = (0, 0, 1)$ where the indices are computed modulo n.

We present a shorter proof of this fact using induction on n.

Claim 7. Proposition 6 is true for n = 1, 2, 3.

Proof. For n = 1, observe that $A(a_1) = (2a_1) = (0)$, and rank $A(0) = \operatorname{rank} A(1) = 0 = \omega(1) = \omega(0)$. For n = 2 we have

$$A(a_1, a_2) = \begin{pmatrix} a_2 & a_1 \\ a_2 & a_1 \end{pmatrix}.$$

It is easily seen that, rank $A(0,0) = 0 = \omega(0,0)$ and rank $A(1,1) = \operatorname{rank} A(1,0) = \operatorname{rank} A(0,1) = 1 = \omega(1,1) = \omega(1,0) = \omega(0,1)$. Let n = 3, in which case

$$A(a_1, a_2, a_3) = \begin{pmatrix} a_2 & a_1 & \\ & a_3 & a_2 \\ a_3 & & a_1 \end{pmatrix}.$$
 (6)

Using the shift-invariance of the rank of A(a), we only need to consider the cases when a equals (0,0,0), (1,0,0), (1,1,0), or (1,1,1). It is easily seen that rank $A(0,0,0) = 0 = \omega(0,0,0)$ and $\omega(1,0,0) = 2 = \operatorname{rank} A(1,0,0)$ and $\omega(1,1,0) = 2 = \operatorname{rank} A(1,1,0)$ and $\omega(1,1,1) = 2 = \operatorname{rank} A(1,1,1)$.

Claim 8. Proposition 6 is true for $a = (0, \ldots, 0)$ and $a = (1, \ldots, 1)$ with $n \ge 3$.

Proof. If a = (0, ..., 0) then A(a) is the zero matrix and rank $A(a) = 0 = \omega(a)$ is clear. If a = (1, ..., 1), then the first n - 1 rows of A(a) are linearly independent, so rank $A(a) \ge n - 1$. On the other hand, (1, ..., 1) is in the kernel of A(a), so rank $A(a) \le n - 1$ and therefore rank $A(a) = n - 1 = \omega(a)$.

We now proceed by induction on n. Let n > 3 be fixed and assume that the claim is true for all vectors $u \in \mathbb{F}_2^k$ with k < n. Let $a \in \mathbb{F}_2^n$. If $a = (0, \ldots, 0)$ or $a = (1, \ldots, 1)$ then the claim is true by Claim 8. Therefore, we may assume that $a \neq (0, \ldots, 0), (1, \ldots, 1)$. Note that from the shift-invariance of χ it follows that the rank of A(a) is invariant under shifts of a. Equivalently, this can also be seen by switching rows and columns. Therefore, we can assume that $a_1 = 1, a_n = 0$. We write the vector a in the following form:

$$a = (\underbrace{1, *, \dots, *, 0}_{=u}, \underbrace{1, \dots, 1}_{=v}, \underbrace{0, \dots, 0}_{=w})$$

More precisely, let k be the last index such that $a_k = 1$ and a_j be the first index such that $a_i = 1$ for all i = j, ..., k. Then $u = (a_1, ..., a_{j-1}) = (1, *, ..., *, 0), v = (a_j, ..., a_k) = (1, ..., 1)$ and $w = (a_{k+1}, ..., a_n) = (0, ..., 0)$. Note that we allow the vector u to be empty. This happens if and only if a = (1, ..., 1, 0, ..., 0), equivalently, j = 1. If a contains at least one occurrence of a 001-pattern, then by shift-invariance we can assume that w contains at least two zeros. Otherwise, w = (0).

Note that wt(a) = wt(u) + wt(v) = wt(u) + (k - j + 1). Now consider the 001-patterns. Any 001-pattern in a either is completely contained inside u, ends exactly at a_j or ends at a_1 . In the first case the 001-pattern is also contained in u. In the second case we know that u = (1, *, ..., *, 0, 0) ends in at least two zeros, and it also has a 001-pattern which ends at a_1 . The last case occurs if and only if w has at least two zeros. It follows that

$$\mathbf{r}(a) = \begin{cases} \mathbf{r}(u) + 1 & w \text{ contains at least two zeros} \\ \mathbf{r}(u) & \text{otherwise.} \end{cases}$$



Then the matrix A(a) has the following form:

Note that A(a) is a block diagonal matrix. The first block is the matrix A(u) with rank $A(u) = \omega(u)$ by the induction hypothesis. This also holds in the degenerate case that u is empty if we then define $\omega(u) = 0$. The second block has rank k - j. Note that if k = j then the second block is empty. The third block has rank 2 if w includes at least two zeros, otherwise it has the form $\begin{pmatrix} 1 \\ 1 \end{pmatrix}$ and has rank 1. Remember that the rank of a block diagonal matrix is the sum of the ranks of the blocks on the diagonal. It follows

$$\operatorname{rank} A(a) = \omega(u) + (k - j) + \begin{cases} 2 & w \text{ contains at least two zeros} \\ 1 & \text{otherwise} \end{cases}$$
$$= \operatorname{wt}(u) + (k - j + 1) + \operatorname{r}(u) + \begin{cases} 1 & w \text{ contains at least two zeros} \\ 0 & \text{otherwise} \end{cases}$$
$$= \operatorname{wt}(a) + \operatorname{r}(a) = \omega(a).$$

For clarity, we also write down how (7) looks in the degenerate cases, namely that u empty, j = k or both. We keep the horizontal and vertical lines to show which blocks vanish. If u is empty and j < k, then a = (1, ..., 1, 0, ..., 0) and

If u is not empty and j = k, then $a = (1, *, \dots, *, 0, 1, 0, \dots, 0)$ and

•

If u is empty and j = k, then $a = (1, 0, \dots, 0)$ and

$$A(a) = \left(\begin{array}{c|cccc} 0 & 1 & & \\ 0 & 0 & & \\ & \ddots & \ddots & \\ & & 0 & 0 \\ 0 & & & 1 \end{array} \right).$$

that

References

- Alex Biryukov, Charles Bouillaguet, and Dmitry Khovratovich. Cryptographic Schemes Based on the ASASA Structure: Black-box, White-box, and Public-key. Cryptology ePrint Archive, Paper 2014/474. 2014. URL: https://eprint.iacr.org/ 2014/474 (visited on 01/18/2024).
- [2] Joan Daemen. "Cipher and hash function design strategies based on linear and differential cryptanalysis". PhD thesis. KU Leuven, 1995.
- [3] Joan Daemen, René Govaerts, and Joos Vandewalle. "An efficient nonlinear shiftinvariant transformation". In: Proceedings of the 15th Symposium on Information Theory in the Benelux. Werkgemeenschap voor Informatie- en Communicatietheorie, 1994.
- [4] Christoph Dobraunig, Maria Eichlseder, Florian Mendel, and Martin Schläffer. "Ascon v1.2: Lightweight Authenticated Encryption and Hashing". In: Journal of Cryptology 34.3 (June 2021), p. 33. ISSN: 1432-1378. DOI: 10/gtfgst.
- [5] Fukang Liu, Santanu Sarkar, Willi Meier, and Takanori Isobe. "The Inverse of χ and Its Applications to Rasta-Like Ciphers". In: Journal of Cryptology 35.4 (Oct. 2022), p. 28. ISSN: 1432-1378. DOI: 10/gtfgn7.
- [6] NIST. SHA-3 Standard: Permutation-Based Hash and Extendable-Output Functions. Tech. rep. Federal Information Processing Standard (FIPS) 202. U.S. Department of Commerce, Aug. 2015. DOI: 10.6028/NIST.FIPS.202.
- [7] Kamil Otal. A Solution to a Conjecture on the Maps $\chi_n^{(k)}$. Cryptology ePrint Archive, Paper 2023/1782. 2023. URL: https://eprint.iacr.org/2023/1782 (visited on 01/18/2024).
- [8] Jan Schoone and Joan Daemen. Algebraic properties of the maps χ_n . Cryptology ePrint Archive, Paper 2023/1708. 2023. URL: https://eprint.iacr.org/2023/1708 (visited on 01/18/2024).