

# Comments on four multi-server authentication protocols using smart card

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## Abstract

Recently, researchers have proposed several nice multi-server authentication protocols. They claim that their protocols are secure and can withstand various attacks. However, after reviewing their schemes, we found that they although are perfect whereas flawed. Due to this observation, in this paper, we list the weakness found in these recent literatures.

**Keywords:** *multi-server, password authentication protocol, smart card, password change, key agreement*

## 1. Introduction

A two-party password authentication protocol for client-server architecture is usually not sufficient when network size increases, especially for an open network such as Internet. Thus, researchers have proposed several multi-server authentication protocols [1-23] to cope with this problem. In 2003, Li *et al.* [5] proposed a multi-server protocol based on the ElGamal digital signature and geometric transformations on an Euclidean plane. Unfortunately, their protocol is broken by Cao and Zhong [8]. In 2004, Juang *et al.* [13] proposed an efficient multi-server scheme and claimed that their scheme is secure, but Chen *et al.* [24] point out it suffers from smart-card-lost attack. In 2004 and 2005, Tsaur *et al.* [3, 4] proposed two multi-server schemes. However, both of their schemes are based on the Lagrange interpolating polynomial which is computationally intensive, and are

broken by Chou *et al.* [16]. In 2006 and 2007, Cao *et al.* [9] and Hu *et al.* [7] each proposed an authentication scheme for a multi-server environment. Both schemes assume that all servers are trustworthy. Nevertheless, this assumption is not always true, as stated in [1]. In 2008, Lee *et al.* [6] proposed an authenticated key agreement scheme for multi-server using mobile equipment. However, their scheme cannot freely add any server. Because when a server is added, all users who want to login to this new server have to re-register at the registration center for obtaining a new smart card. This increases the registration center's card-issue cost. Also, in 2008, Tsai [1] proposed an efficient multi-server authentication scheme and claimed that his protocol can withstand seven known attacks. Yet, Chen *et al.* [15] and Wang *et al.* [17] finds it is vulnerable to the server spoofing attack. In addition, Tsaura *et al.* [21] also find it is vulnerable to the man-in-the-middle attack. In 2009, Liao *et al.* [2] proposed a secure dynamic ID scheme for multi-server environment. They claimed their protocol is secure. However, Chen *et al.* [15] and Hsiang and Shih [14] both find it suffers server spoofing attack by an insider server. Further, [14] propose an improvement on the protocol. In recent two years, literatures [19-22] each not only point out the security flaws in previous schemes (as stated above) but also propose a secure protocol in multi-server environment. Lee *et al.* [19] and Sood *et al.* [20] both find the improvement in [14] is still insecure. In 2011, Tsaura *et al.* [21] pointed out [13] have man-in-the-middle attack, and Li *et al.* [22] indicated that [20] has leak-of-verifier attack and stolen smart card attack in 2012. Finally also in 2012, Hwang *et al.* [18] propose an improved multi-server authentication protocol based on bilinear pairings, and Liao *et al.* [23] propose a novel multi-server authentication scheme for mobile clients. However, after examining schemes [19-23], this study found there still existing some deficiency in each proposal. In this paper, we will first show the deficiencies found, then show our scheme and examine its security.

The rest of this paper is organized as follows: Section 2 demonstrates the vulnerabilities existing in schemes [19-23]. Section 3 demonstrates a novel protocol and Section 4 analyzes its security. The discussions and comparisons are made in Section 5. Finally, a conclusion is given in Section 6.

## 2. Deficiencies in literatures [19-23]

In this section, we demonstrate the deficiency in each scheme.

- Lee *et al.*'s protocol [19]

Lee *et al.*'s protocol is a secure dynamic ID based scheme, but might suffer insider attack. Because any insider server has the secrecy  $h(x || y)$  and  $h(y)$ , a malicious insider server  $S_m$  can intercept the message  $\{CID_i, P_{ij}, Q_i, N_i\}$  sent from the user, and

calculate  $T_i = P_{ij} \oplus h(h(y) \parallel N_i \parallel \text{SID}_j)$  and  $A_i = h(T_i \parallel h(y) \parallel N_i)$ . Consequently,  $S_m$  can obtain  $h(b \oplus \text{PW}_i)$  by computing  $\text{CID}_i \oplus h(T_i \parallel A_i \parallel N_i)$ , and then computes  $B_i = h(h(b \oplus \text{PW}_i) \parallel h(x \parallel y))$ . Then,  $S_m$  can masquerade as the user to login any server by generating a random  $N_i'$  and computing the corresponding login message.

- Sood *et al.*'s protocol [20]

Li *et al.* [22] indicate that [20] suffers leak-of-verifier attack, stolen smart card attack, and has incorrect authentication and session key agreement phase.

- Tsaaur *et al.*'s protocol [21]

Tsaaur *et al.*'s protocol suffers the smart card lost password guessing attack. Because if the attacker intercepted  $M_{ij} = \{ E\_T_{ij}, A_{ij}, \text{UID}_i, \text{Ev}_{ij}(\text{ru}_k, h(\text{UID}_i)) \}$ , then from the stored values  $\{\text{UID}_i, \mu_i, E\_T_{ij}, \text{and } A_{ij}\}$  in the smart card, he can guess  $\text{UID}_i$ 's password  $\text{PW}_i$  as  $\text{gpw}$  and compute  $h(\text{gpw})$ . He computes  $h\text{UID}' = h(\text{UID}_i)$ , and calculates  $v_i' = \mu_i \oplus h(\text{gpw})$  and  $v_{ij}' = h(v_i', \text{SID}_j)$  to decrypt  $\text{Ev}_{ij}(\text{ru}_k, h(\text{UID}_i))$  in  $M_{ij}$ . He verifies whether  $h(\text{UID}_i)$  in the decryption result is equal to the computed  $h\text{UID}'$ . If it is, the attacker can confirm that the user's password is  $\text{gpw}$ .

- Li *et al.*'s protocol [22]

Li *et al.*'s protocol is efficient and secure, but it might suffer impersonation attack. We describe the reasons as follows:

From the values  $D_i, b, h(y), C_i, D_i,$  and  $E_i$  stored in the smart card, a user  $\text{ID}_i$  can obtain  $h(y \parallel x)$  by first computing  $A_i = h(b \parallel P_i)$  and  $B_i = D_i \oplus h(\text{ID}_i \parallel A_i)$ , where  $P_i$  is  $\text{ID}_i$ 's password. Then, after  $S_j$  has sent login message  $(F_i, G_i, P_{ij}, \text{CID}_i, \text{SID}_j, K_i, M_i)$  to CS,  $\text{ID}_i$  can obtain  $h(y \parallel x)$  by computing  $B_i \oplus P_{ij} \oplus h(h(y) \parallel N_{i1} \parallel \text{SID}_j)$ , where  $N_{i1}$  is a random number generated by the smart card and  $P_{ij} = E_i \oplus h(h(y) \parallel N_{i1} \parallel \text{SID}_j)$ . On the other hand, when acting as an attacker, from the login request message sent from  $S_j$  to the CS, an insider having  $h(y)$  stored in the smart card can compute  $\text{ID}_i$ 's relevant parameters  $N_{i1} = h(y) \oplus F_i$ ,  $E_i = P_{ij} \oplus h(h(y) \parallel N_{i1} \parallel \text{SID}_j)$ ,  $B_i = E_i \oplus h(y \parallel x)$ , and  $A_i = \text{CID}_i \oplus h(B_i \parallel F_i \parallel N_{i1})$ . After obtaining  $A_i, B_i, E_i,$  and  $N_{i1}$ , the insider computes  $h(A_i \parallel B_i \parallel N_{i1})$  and  $(N_{i2} \oplus N_{i3}) = T_i \oplus h(A_i \parallel B_i \parallel N_{i1})$ , where  $T_i = N_{i2} \oplus N_{i3} \oplus h(A_i \parallel B_i \parallel N_{i1})$  is in the submitted message from CS to server  $S_j$ . Then, the insider can compute the session key  $\text{sk} = h(h(A_i \parallel B_i) \parallel (N_{i1} \oplus N_{i2} \oplus N_{i3}))$ . Moreover, for possessing  $A_i, B_i, E_i,$  the insider can generate a random  $N_{i1}'$ , compute  $F_i = h(y) \oplus N_{i1}'$ ,  $G_i = h(B_i \parallel A_i \parallel N_{i1}')$ ,  $P_{ij} = E_i \oplus h(h(y) \parallel N_{i1}' \parallel \text{SID}_j)$ , and  $\text{CID}_i = A_i \oplus h(B_i \parallel F_i \parallel N_{i1}')$ . That is, even though the insider does not know whom  $\text{CID}_i$  stands for, the insider can

impersonate a legal user to login to the server successfully. Besides, it requires that the CS must be kept on-line. This might cause the CS to be a bottleneck in the system.

- Liao *et al.*'s protocol [23]

Liao et al. proposed a perfect multi-server scheme for mobile clients, but a deficiency of suffering password guessing attack may exist. Because, once the smart card of a user is lost, an insider server, who knows the user's  $ID_i$  and  $T_i$  (in the ID table sent from the register server RS to all the service servers periodically over a secure channel), the value set  $\{ID_i, RegID_i, b_i\}$  stored in the smart card, and has ever intercepted the user's login message  $(ID_i, M_i, B_{ij}, R_i)$  to the service server  $SS_j$ , can launch a password guessing attack. We depict the scenario as follows. The insider server guesses a password  $gpw$ , computes  $DID_i = RegID_i \oplus h(gpw \parallel b_i) \cdot Pub_{RS}$ ,  $d_{ij} = h(ID_i \parallel SID_j \parallel M_i \parallel R_i)$ , and  $B_{ij} - d_{ij} \cdot DID_i = r_i \cdot DID_i$ . He then can confirm the correctness of  $gpw$  by verifying whether equation  $e(r_i \cdot DID_i, P) = e(M_i, T_i \cdot Pub_{RS})$  holds. If it does, the insider server can assure that the password  $gpw$  he guesses is right. We deduce the equation as follows.

$$e(r_i \cdot DID_i, P) = e(r_i \cdot T_i \cdot s_{RS} \cdot QID_i, P) = e(r_i \cdot QID_i, T_i \cdot s_{RS} \cdot P) = e(M_i, T_i \cdot Pub_{RS}),$$

where  $DID_i = (T_i \cdot s_{RS}) \cdot QID_i$  is computed by RS in user registration phase.

### 3. Conclusion

In this article, we have listed the weakness in recent four studies in multi-server environment. The analyses show that these schemes are deficient and need further improvement.

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